# Social Choice — Finite Sets

# 1 Individual Preferences: Assumptions

- 1. A finite set N of individuals: i = 1, 2, ..., n, with  $n \ge 2$
- 2. A finite set X of alternatives:  $|X| \ge 3$
- 3. Individual preference  $R^i \in \Psi$ : a weak order on X $\Box \Psi \equiv$  Set of all possible weak orders on X (complete, reflexive, transitive)
- 4. Preference profile  $\rho = (R^1, \dots, R^n) \in \Psi^n$ : n-tuple of weak orders.  $\Box \Psi^n \equiv \text{Set of all possible preference profiles.}$
- 5. Restricted preference profile:  $\rho|_S = (R^1|_S, \cdots, R^n|_S), S \subseteq X$
- 6. Collective decision rule (CDR) f(ρ), denoted R, is a mapping: Ψ<sup>n</sup> → ℜ
  □ ℜ ≡ Set of all *complete* binary relations.
  - $\triangleright$  f is a preference aggregation rule, yielding a social ranking.
  - $\triangleright$  Also called social welfare function (SWF) or constitution [Green/Laffont 1979].
- 7. Some sets of voters:
  - $R(x, y; \rho) \equiv \{i \in N \mid {}_x R^i{}_y\}$
  - $P(x, y; \rho) \equiv \{i \in N \mid {}_{x}P^{i}{}_{y}\}$
  - $I(x,y;\rho) \equiv \{i \in N \mid {}_xI^i{}_y\}$

Note:  $R(x, y; \rho) \neq R(y, x; \rho)$ , and  $P(x, y; \rho) \neq P(y, x; \rho)$ 

### 2 Collective Decision Rules

- 1. Some examples:
  - Constant rule:  $f(\rho) = f(\rho'), \ \forall \rho, \rho' \in \Psi^n$
  - Pareto/unanimity rule:  ${}_{x}P_{y}$  iff  ${}_{x}P^{i}{}_{y}, \forall i \in N$
  - Majority rule:  $_{x}P_{y}$  iff  $|P(x,y;\rho)| > \frac{n}{2}; \ _{x}R_{y}$  iff  $|R(x,y;\rho)| \ge \frac{n}{2}$
  - Plurality rule:  $_xP_y$  iff  $|P(x,y;\rho)| > |P(y,x;\rho)|$ ;  $_xR_y$  iff  $|R(x,y;\rho)| \ge |R(y,x;\rho)|$
  - Border rule:  $_xP_y$  iff  $\sum_i r_i(x) < \sum_i r_i(y)$ ;  $r_i(\cdot) \equiv$  assigned rank number
- 2. Properties of CDR  $f(\rho)$ :
  - Universal Domain (UD): f has full  $\Psi^n$  as its domain.

- Non-dictatorial (ND):  $\not\exists i \in N$  s.t.  $\forall \rho \in \Psi^n, \forall x, y \in X : {_xP^i}_y \Rightarrow {_xP_y}$
- Weakly Paretian (WP):  $\forall \rho \in \Psi^n, \forall x, y \in X : {_xP^i}_y, \forall i \in N \Rightarrow {_xP_y}$
- Independence of irrelevant alternatives (IIA):  $\forall \rho, \rho' \in \Psi^n, \forall x, y \in X$ :

$$\rho|_{\{x,y\}} = \rho'|_{\{x,y\}} \quad \Rightarrow \quad f(\rho)|_{\{x,y\}} = f(\rho')|_{\{x,y\}}$$

- 3. Some examples:
  - Constant rule: ND, IIA, but not WP.
  - Dictatorship: WP, IIA, but not ND.
  - Pareto rule: ND, WP, IIA, but not COMP, not TRAN (only Q-TRAN).
  - Borda rule:<sup>1</sup> ND, WP, but not IIA.
  - Majority rule: ND, WP, IIA, but not ACYC (hence not TRAN).

### **3** Rationality of Collective Decision Rules

**Def 1:** A CDR  $f(\rho)$  is said to be ACYC/Q-TRAN/TRAN if it is so for any profile  $\rho \in \Psi^n$ .

**Def 2 (Decisiveness)** Coalition  $L \subseteq N$  is:

• semi-decisive over (x, y):  ${}_x \tilde{D}^L{}_y$  if

$$\forall \rho: (\forall i \in L, {_xP^i}_y) \& (\forall j \notin L, {_yP^j}_x) \Rightarrow {_xP_y}$$

• decisive over (x, y):  ${}_{x}D^{L}{}_{y}$  if

$$\forall \rho: \forall i \in L, \, _{x}P^{i}{}_{y} \Rightarrow _{x}P_{y}$$

• decisive if L is decisive over any pair  $(x, y) \in X^2$ 

**Lmm:** If f is Q-TRAN, WP, and IIA, then for any coalition  $L \in N$ :

(1)  $\exists x, y \in X, \ _x \tilde{D}^L{}_y \Rightarrow \forall z_{(\neq x,y)} \in X : \ _x D^L{}_z \text{ and } y D^L{}_z$ (2)  $\exists x, y \in X, \ _x \tilde{D}^L{}_y \Rightarrow \forall r, s \in X : \ _r D^L{}_s \blacksquare$ 

Thm (Arrow 1951) For  $|X| \ge 3$  and  $n \ge 2$ :

If f is UD, TRAN, WP, and IIA, then it must be dictatorial.

#### Def 3: Veto power

- Agent  $i \in N$  has a veto for (x, y) if  $\forall \rho \in \Psi^n, \ _x P^i{}_y \Rightarrow \sim _y P_x$
- Agent  $i \in N$  has a veto if i has a veto for all (x, y).
- f is oligarchic if  $\exists L (\subseteq N)$  that is decisive, and every  $i \in L$  has a veto.

Thm (Gibbard 1973) If f is UD, Q-TRAN, WP, and IIA, then it is oligarchic.

<sup>&</sup>lt;sup>1</sup>Or any point voting system.

**Def 4:** Winning/decisive coalition set  $\mathcal{L}(f) \equiv \{L \subseteq N \mid L \text{ is decisive under rule } f\} \subseteq 2^N$   $\triangleright$  If  $L_1, L_2 \in \mathcal{L}(f)$ , then  $L_1 \cap L_2 \neq \emptyset$ . [Otherwise conflict may result.]  $\triangleright$  If f is WP, then  $N \in \mathcal{L}(f)$ .

**Def 5:** Winning coalition set  $\mathcal{L}(f)$  is:

- monotonic:  $L \in \mathcal{L}(f)$  and  $L \subseteq L' \Rightarrow L' \in \mathcal{L}(f)$
- proper:  $L \in \mathcal{L}(f) \Rightarrow N \setminus L \notin \mathcal{L}(f)$

**Lmm:** For any f, coalition set  $\mathcal{L}(f)$  is monotonic and proper.

**Def 6 (Collegial)** f is collegial (決策核心制) if

$$\bigcap_{L \in \mathcal{L}(f)} L \neq \emptyset$$

- ▷ Collegium (核心成員) of a collegial f:  $K(f) \equiv \bigcap_{L \in \mathcal{L}(f)} L$
- $\triangleright$  The collegium is necessary for any decision, but may not be sufficient.

(eg) Decision rule of the UN Security Council

Thm (Brown 1975) If  $|X| \ge n$  and f is ACYC and WP, then it is collegial.

**Ex:** Decision rule of UN Security Council is collegial, not oligarchic.<sup>2</sup>  $\Box$ 

**Thm:** If |X| > n and f is ACYC, WP, and IIA, then  $\exists i \in N$  with a veto over some (x, y).

**Def 7:** Derived rule  $f_{\mathcal{L}}$  for  $\mathcal{L} \subseteq 2^N$ :  ${}_xP^{\mathcal{L}}{}_y \rightleftharpoons \exists L \in \mathcal{L}$  s.t.  ${}_xP^i{}_y, \forall i \in L$ 

 $\triangleright f$  is simple if  $f = f_{\mathcal{L}(f)}$ .

 $\rhd f$  is more resolute than  $f_{\mathcal{L}(f)}$ , since  $_x P^{\mathcal{L}(f)}_y \Rightarrow _x P_y$ .

**Def 8:** For all  $\rho, \rho' \in \Psi^n$  and  $x, y, a, b \in X$ , a rule f is:

- decisive iff  $[P(x, y; \rho) = P(x, y; \rho') \& {}_{x}P_{y}] \Rightarrow {}_{x}P'_{y}$
- neutral iff  $[P(x,y;\rho) = P(a,b;\rho') \& P(y,x;\rho) = P(b,a;\rho')] \Rightarrow [_xP_y \rightleftharpoons _aP'_b]$
- monotonic iff  $[P(x,y;\rho) \subseteq P(x,y;\rho') \& R(y,x;\rho) \subseteq R(x,y;\rho') \& {}_{x}P_{y}] \Rightarrow {}_{x}P'_{y}$

 $\triangleright$  Neutrality implies that names of agents do not matter.

- **Thm:** f is simple iff f is decisive, neutral, and monotonic.
  - $\rhd$  Plurality rule is not decisive, hence not simple.

**Def 9:** A simple rule f is a q-rule  $(q > \frac{n}{2})$  iff  $\mathcal{L}(f) = \{L \subseteq N : |L| \ge q\}$ .

- $\triangleright$  Pareto rule is a *q*-rule with q = n.
- $\triangleright$  Majority rule is a q-rule with  $q = \frac{n+1}{2}$ .

 $<sup>^{2}</sup>$ It is ACYC, not Q-TRAN.

**Def 10:** Nakamura number for rule *f*:

$$v(f) \equiv \min\{ |\mathcal{L}'| : \mathcal{L}' \subseteq \mathcal{L}(f), \cap_{L \in \mathcal{L}'} L = \emptyset \}$$

**Ex:** For majority rule:

• n = 3: v(f) = 3,  $\mathcal{L}(f) = \{\underline{\{1,2\}}, \underline{\{1,3\}}, \underline{\{2,3\}}, \{1,2,3\}\}$ . • n = 4: v(f) = 4,  $\mathcal{L}(f) = \{\underline{\{1,2,3\}}, \underline{\{1,2,4\}}, \underline{\{1,3,4\}}, \underline{\{2,3,4\}}, \{1,2,3,4\}\}$ .  $\Box$ 

**Lmm:** For any rule f,  $v(f) \ge 3.^3$ 

**Lmm:** If f is not collegial, then  $v(f) \leq n$ .

- Thm (Nakamura) A simple rule f is ACYC iff |X| < v(f).
- **Def 11:** A simple rule f is strong iff  $\forall L: L \notin \mathcal{L}(f) \Rightarrow N \setminus L \in \mathcal{L}(f)$ .  $\triangleright$  Majority rule (with n odd, or n even with a tie-breaker) is strong.

**Lmm:** If f is collegial and strong, then it is dictatorial.

**Lmm:** If f is non-collegial and strong, then v(f) = 3.

**Lmm:** If f is a q-rule with  $q \le n$ , then  $v(f) = \frac{n}{n-q}$ .

**Ex:** For majority rule f: v(f) = 3 (if  $n \neq 4$ ) and v(f) = 4 (if n = 4).  $\Box$ 

**Cor:** A non-collegial, strong simple rule is ACYC iff |X| < 3.

**Cor:** A non-collegial q-rule is ACYC iff  $|X| < \frac{n}{n-q}$ .  $\triangleright$  Majority rule with  $|X| \ge 3$  (when  $n \ne 4$ ) is not acyclic.

Def 12: Blocking coalition set:

$$\mathcal{B}(f) \equiv \{ L \subseteq N \,|\, N \setminus L \notin \mathcal{L}(f) \}$$

Losing coalition set:

$$\mathcal{S}(f) \equiv 2^N - B(f) = \{L \subseteq N \mid N \setminus L \in \mathcal{L}(f)\}\$$

Non-winning blocking coalitions:

$$\{L \subseteq N \mid L \in B(f), \ L \notin \mathcal{L}(f)\}$$

 $\triangleright$  For strong  $f, B(f) = \mathcal{L}(f)$ : a coalition either wins or loses.

<sup>&</sup>lt;sup>3</sup>Two disjoint coalitions cannot both be decisive.

**Ex:** q-rule with n = 100 and q = 60:

 $\mathcal{L}(f) = \{L \subseteq N : |L| \ge 60\}$  $\mathcal{B}(f) = \{L \subseteq N : |L| > 40\}$ 

 $\mathcal{S}(f) = \{L \subseteq N : |L| \le 40\} \square$ 

Remark (Fundamental Dilemma) Trade-off among ACYC, Equality, and Resoluteness.

- (1) Pareto rule: most EQU, most ACYC, but worst in RES.
- (2) Majority rule: most EQU, most RES, but worst in ACYC.
- (3) Dictatorship: most ACYC, most RES, but worst in EQU.

**Def 13 (Core)**  $C_f(\rho, X) \equiv M(f(\rho), X)$ 

- Core of a simple rule f is not empty iff |X| < v(f) [Nakamura Thm]
- Theory predicts only in very restricted situations.
- Core is typically empty!

# 4 Social Decision Functions (SDF)

**Def 14 (Social decision function)** An SDF  $\delta(\rho)$  is a mapping:  $\Psi^n \mapsto X$ .

- $\triangleright$  SDFs assign to each preference profile  $\rho$  an element in X.
- $\triangleright$  Also known as social choice function (SCF) [Green/Laffont 1979].

**Def 15:** SDF  $\delta(\rho)$  is manipulable at  $\rho = (R^1, \dots, R^n)$  if there exists  $R^{i'} \in \Psi$  such that:

$$[\delta(R^1,\dots,R^{i\prime},\dots,R^n)]P^i[\delta(R^1,\dots,R^i,\dots,R^n)]$$

**Def 16:** SDF  $\delta(\rho)$  is strongly individually incentive compatible (SIIC) if it is not manipulable at any preference profile  $\rho$ .

 $\triangleright$  Truthful revelation of preferences is dominant strategy for all *i* and for any  $\rho$ :

$$\forall \rho \in \Psi^n, \ \forall R^{i'} \in \Psi: \ _{[\delta(R^1, \cdots, R^i, \cdots, R^n)]} R^i_{[\delta(R^1, \cdots, R^{i'}, \cdots, R^n)]}$$

 $\triangleright$  Also called cheat-proof, strategy-proof, or straightforward. <u>NB</u>: SIIC is weaker then *Group-/coalition-nonmanipulability*.

**Def 17:** If SDF  $\delta(\cdot)$  has range  $S \subseteq X$ , then it is called SDF with range S.

**Def 18:** SDF  $\delta$  with range S is dictatorial if  $\exists i \in N$  (the dictator) such that:

$$\forall \rho \in \Psi^n, \forall x \in S : x \neq \delta(\rho) \Rightarrow \delta(\rho) R^i_x$$

 $\triangleright$  There is an agent whose favorite alternative is always the social choice.

**Lmm** If there exists  $i \in N$  such that

 $\delta(R^1, \cdots, R^i, \cdots, R^n) = x, \ \delta(R^1, \cdots, R^{i'}, \cdots, R^n) = y; \ x \neq y$ 

and if  ${}_{x}P^{i}{}_{y}$  and  ${}_{x}P^{i}{}_{y}'$ , then  $\rho$  is manipulable at either  $(R^{1}, \dots, R^{i}, \dots, R^{n})$  or  $(R^{1}, \dots, R^{i'}, \dots, R^{n})$ .

(Pf) Obvious by the definition of manipulability.  $\Box$ 

**Lmm** Let  $\delta$  be SIIC with range  $S \subseteq X$ . If  $T \subseteq S$  and  $\rho = (R^1, \dots, R^n)$  is a profile such that

 $\forall i \in N, \forall x, y \in S \text{ with } x \in T, y \notin T : x P^{i}_{y},$ 

then  $\delta(\rho) \in T$ .

Thm (Gibbard 1973/Satterthwaite 1975) If  $|S| \ge 3$ , then any SDF with range S satisfying SIIC and UD is dictatorial.

 $\triangleright$  Allowing more complex strategy space does not help!

Generalization Follow-up research:

(1) Domain restriction [Maskin; Kalai/Muller]

- (2) Imposed Structure: free disposal, neutral agent
- (3) Statistical info about taste distribution [Grandmont]
- (4) Random social lottery [Gibbard]