

私人捐獻均衡

1. The Cornes-Sandler Framework

1.1. Additive Public Good

- n consumers
- 2 goods:
 - x : private good (price = 1)
 - Q : public good (price = p)

- Private contribution to PG:

$$q_i$$

- Additive PG:

$$Q = \sum q_i, \quad \tilde{Q}_i \equiv \sum_{j \neq i} q_j$$

- Pure altruistic preference:

$$U_i(x_i, Q) = U_i(x_i, q_i + \tilde{Q}_i)$$

▷ U is strictly quasi-concave

- Individual budget allocation:

$$x_i + pq_i = I_i$$

- Utility-max:

$$\max_{x, q} U(x, q + \tilde{Q}) \quad \text{s.t.} \quad x + pq = I \quad (1)$$

or:

$$\max_q V(q, \tilde{Q}) \equiv U(I - pq, q + \tilde{Q}) \quad (2)$$

- Iso-utility (IU) curve:

$$V(q, \tilde{Q}) = k$$

▷ Slope:

$$\frac{d\tilde{Q}}{dq} = \frac{-\partial V/\partial q}{\partial V/\partial \tilde{Q}} = \frac{pU_x - U_Q}{U_Q} = p\frac{U_x}{U_Q} - 1 \quad (3)$$

E Show that “better sets”

$$B \equiv \{(q, \tilde{Q}) \mid V(q, \tilde{Q}) \geq k; k \in R\}$$

of function $V(q, \tilde{Q})$ are convex. \square

- Nash reaction function: Fig. 1

$$q(\tilde{Q})$$

▷ Nash reaction curve NN :

$$\frac{d\tilde{Q}}{dq} = p\frac{U_x}{U_Q} - 1 = 0$$

hence:

$$\frac{U_Q}{U_x} = p \quad (4)$$

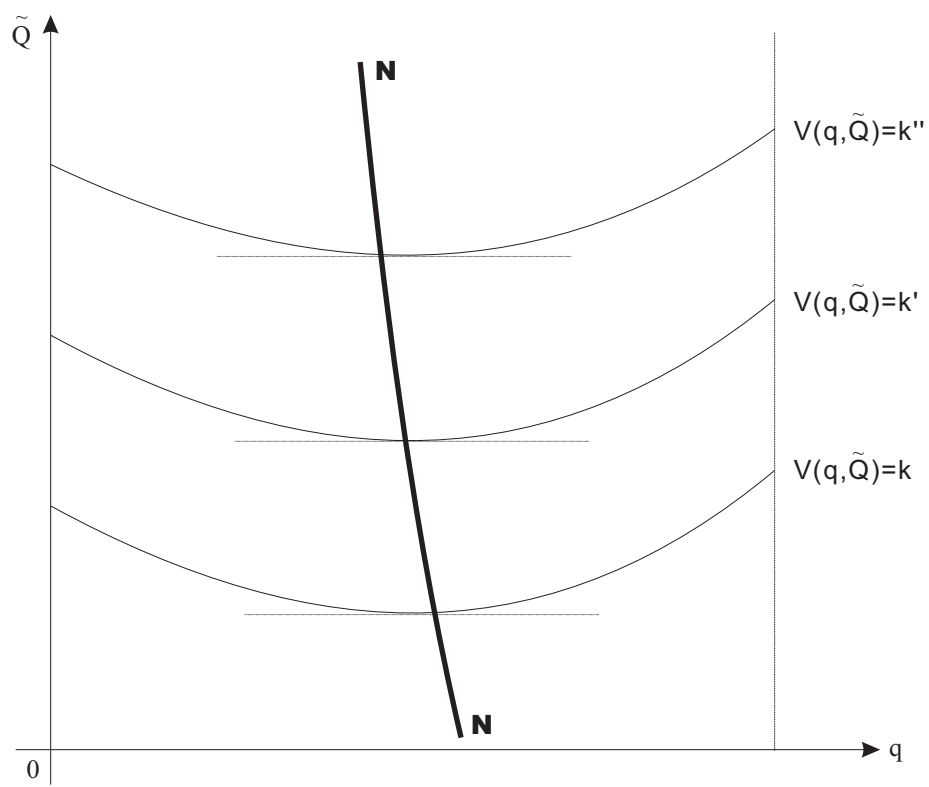
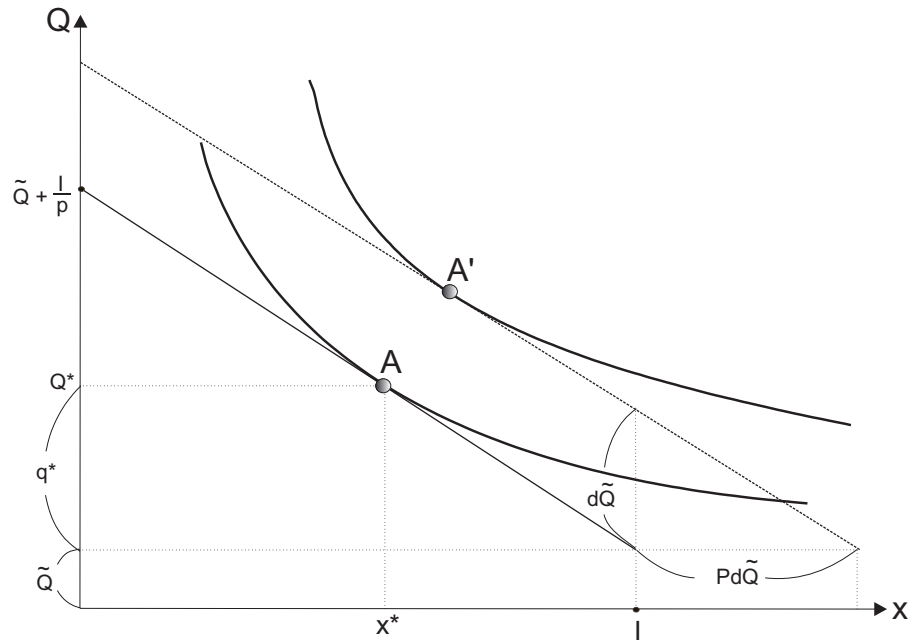


Figure 1: Iso-utility curves in q - \tilde{Q} plane

1.2. Individual Nash Choice

- Choice of (x, q) :

$$\begin{aligned}\frac{dQ}{d\tilde{Q}} &= 1 + \frac{dq}{d\tilde{Q}} \\ \frac{dq}{d\tilde{Q}} &= \frac{dQ}{d\tilde{Q}} - 1\end{aligned}\quad (5)$$



- Choice of (x, Q) : substitute $[q = Q - \tilde{Q}]$ into (1)

$$\max_{x, Q} U(x, Q) \quad \text{s.t.} \quad x + pQ = I + p\tilde{Q} \quad (6)$$

hence:

$$\frac{dQ}{d\tilde{Q}} = p \frac{dQ}{dI} \quad (7)$$

or, by substituting (7) into (5):

$$\frac{dq}{d\tilde{Q}} = p \frac{dQ}{dI} - 1 \quad (8)$$

- Slope of NN :

- x and Q are both normal:

- ▷ A' to the up-right of A

$$1 > p \frac{dQ}{dI} > 0$$

$$0 > \frac{dq}{d\tilde{Q}} > -1, \quad \frac{d\tilde{Q}}{dq} < -1 \quad (9)$$

- ▷ NN has negative slope (< -1)

- x is normal, but Q is inferior:

- ▷ A' to the lower-right of A

$$p \frac{dq}{dI} < 0$$

$$\frac{dq}{d\tilde{Q}} < -1, \quad -1 < \frac{d\tilde{Q}}{dq} < 0 \quad (10)$$

- ▷ NN has negative slope (> -1)

- Q is normal, but x is inferior:

- ▷ A' to the up-left of A

$$p \frac{dQ}{dI} > 1$$

$$\frac{dq}{d\tilde{Q}} > 0, \quad \frac{d\tilde{Q}}{dq} > 0 \quad (11)$$

- ▷ NN has positive slope

2. 2-player Normal-good Equilibrium

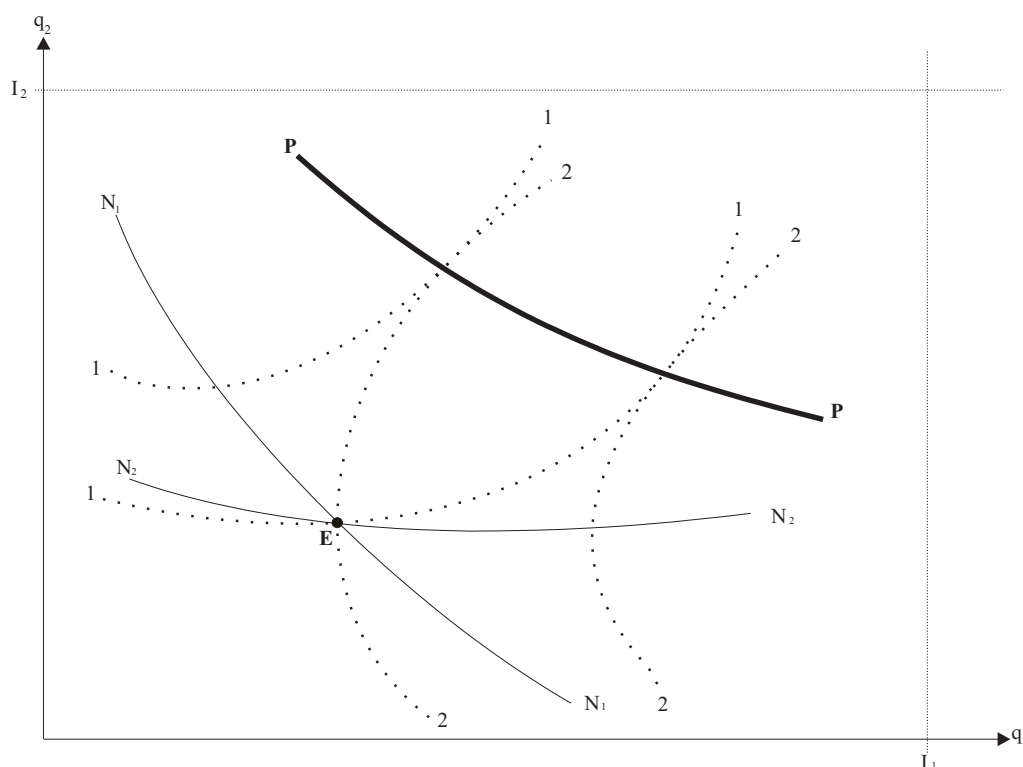
- Individual util-max:

$$V^1(\bar{q}_2 | I_1) \equiv \max_{x_1, q_1} U(x_1, q_1 + \bar{q}_2) \quad \text{s.t.} \quad x_1 + pq_1 = I_1$$

$$V^2(\bar{q}_1 | I_2) \equiv \max_{x_2, q_2} U(x_2, \bar{q}_1 + q_2) \quad \text{s.t.} \quad x_2 + pq_2 = I_2$$

- Slope of Nash reaction curves:

$$\frac{dq_1(q_2)}{dq_2} < -1, \quad \frac{dq_2(q_1)}{dq_1} < -1$$



- Nash equilibrium: intersection of Nash curves NN_1 and NN_2

- Pareto efficiency: tangency of iso-utility curves (by slope (3))

$$\frac{dq_2}{dq_1} = p \frac{U_x^1}{U_Q^1} - 1$$

$$\frac{dq_1}{dq_2} = p \frac{U_x^2}{U_Q^2} - 1$$

▷ Tangency on (q_1-q_2) plane:

$$p \frac{U_x^1}{U_Q^1} - 1 = \left[p \frac{U_x^2}{U_Q^2} - 1 \right]^{-1} \quad (12)$$

▷ Hence the Samuelson FOC:

$$\text{MRS}_1 + \text{MRS}_2 \equiv \frac{U_Q^1}{U_x^1} + \frac{U_Q^2}{U_x^2} = p$$

- Under-provision of Nash:

$$Q^E = q_1 + q_2 < PP$$

- Interpretation: Prisoners' Dilemma
- Nash must exist, but may not be unique!

3. Free-riding

- Definition:

1. Micro-level:

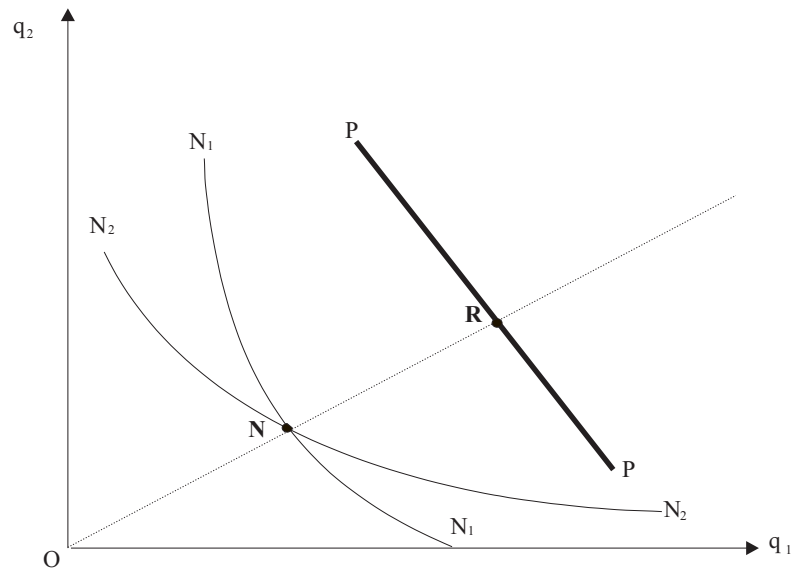
$$q'(\tilde{Q}) < 0$$

2. Systemic level:

$$Q^E = q_1 + q_2 < PP$$

- Index/measurement:

$$\delta \equiv \frac{ON}{OR} \in (0, 1)$$



- n -player case: on q - \tilde{Q} space

- Symmetric solution: both Nash and PO have slope $(n - 1)$:

$$q = \frac{Q}{n} = \frac{\tilde{Q}}{n - 1}$$

- Symmetric Nash: intersection of NN and ray of slope $(n - 1)$

– Symmetric Pareto: tangency of EU and ray of slope $(n - 1)$

□ Since

$$p \frac{U_x}{U_Q} - 1 = n - 1$$

we have the Samuelson foc:

$$p = n \cdot \text{MRS}^{Q,x} \quad \square$$

• Population effect: as $n \uparrow$

– Quasi-linear:

$$U(x, Q) = x + f(Q), \quad f' > 0, \quad f'' < 0$$

▷ δ goes down as $n \uparrow$

□ Since $dQ/dI = 0$,

$$\frac{dq}{d\hat{Q}} = p \frac{dQ}{dI} - 1 = -1$$

Nash $\hat{Q} = q + \tilde{Q}$ is fixed. However, Pareto Q^* must satisfy

$$nf'(Q^*) = p$$

and hence will go up with n . ■

– Cobb-Douglas (or log-linear):

$$U(x, Q) = \log x + r \log Q$$

▷ δ goes down as $n \uparrow$, eventually approaching 0.

□ Now Nash is:

$$\hat{Q} = \frac{r}{n+r} \sum_i I_i$$

while Pareto is:

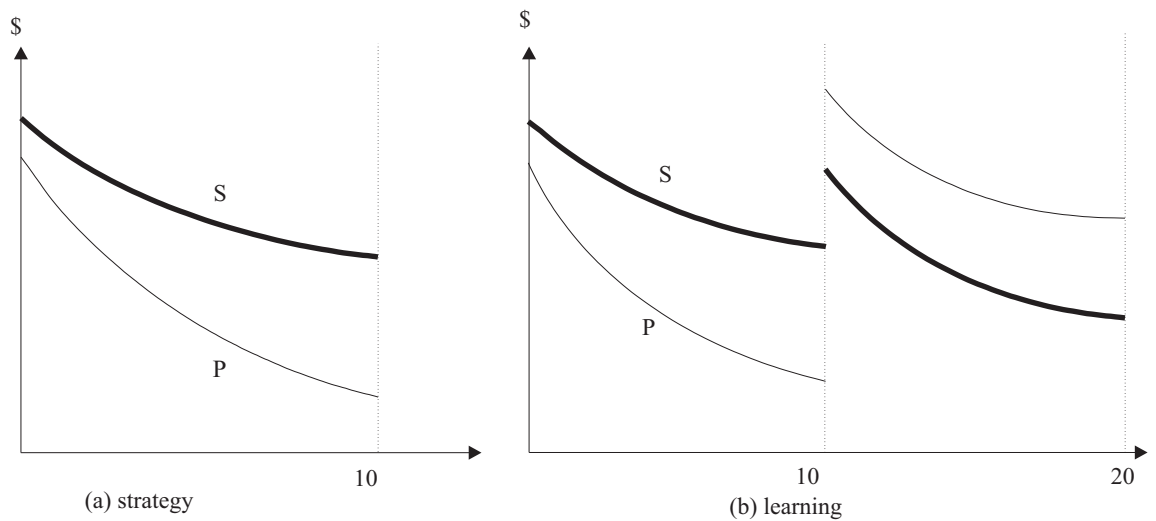
$$Q^* = \frac{r}{1+r} \sum_i I_i$$

So δ goes to 0 as $n \uparrow$. See Laffont, pp.39–41. ■

- Data testing: Haan-Kooreman [JPuE 2002/v83, pp. 277–91]
 - ▷ Weak evidence that free riding increases with group size.

- Experiment design:

- Andreoni [JPuE, 1988, 37:291–304]
- Learning v. Strategy



- Restart effect: Ambrus and Pathak [JPuE, 2011/v95, pp. 500–512]
 - * Two player types: *selfish* v. *reciprocal*

4. The Neutrality/Invariance of Income Redistribution

- Bergstrom et al. [JPuE 1986]

$$U_i(x_i, G), \quad G = \sum_i g_i$$

- Original (unique) Nash equilibrium:

$$(x_1^*, \dots, x_n^*, g_1^*, \dots, g_n^*, G^*)$$

- Income redistribution among contributors ($g_i > 0$):

$$w'_i = w_i + \Delta w_i, \quad \sum_i \Delta w_i = 0$$

▷ Amount distributed not too big:

$$|\Delta w_i| \leq g_i^*, \quad \text{if } \Delta w_i < 0$$

$$\Delta w_i \leq G_{-i}^*, \quad \text{if } \Delta w_i > 0$$

- New (unique) Nash equilibrium: Fig. 2

$$(x'_1, \dots, x'_n, g'_1, \dots, g'_n, G')$$

with:

$$x'_i = x_i^*$$

$$g'_i = g_i^* + \Delta w_i$$

$$G' = G^*$$

$$U_i(x'_i, G') = U_i(x_i^*, G^*)$$

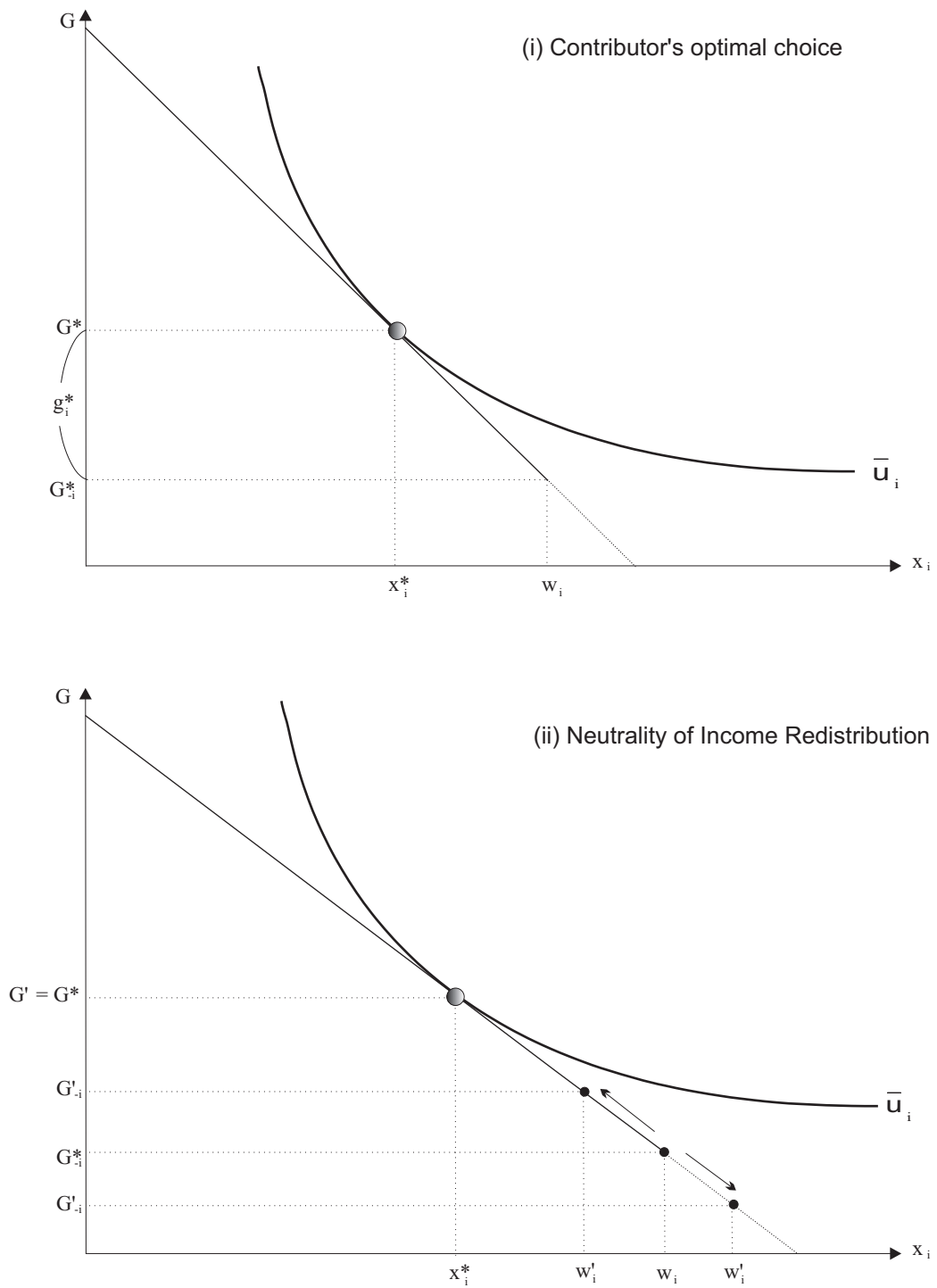


Figure 2: The Neutrality Hypothesis

5. Neutrality of Forced Contributions

5.1. Warr [JPuE 1982]

- Three consumers: rich are altruistic

$$U^1 = u^1(c^1, c^3), \quad U^2 = u^2(c^2, c^3)$$

$$U^3 = u^3(c^3)$$

- Voluntary transfer:

$$c^i = I^i - v^i, \quad i = 1, 2$$

$$c^3 = I^3 + [v^1 + v^2]$$

- Original Nash equilibrium:

$$\max_{v^1} u^1(c^1, c^3) \quad \text{s.t.} \quad c^1 = I^1 - v^1$$

$$\max_{v^2} u^2(c^2, c^3) \quad \text{s.t.} \quad c^2 = I^2 - v^2$$

▷ Interior FOC:

$$u_i^i(c^i, c^3) = u_3^i(c^i, c^3); \quad i = 1, 2$$

- NE is not PO: with 1 and 2 both contributing \$1 more

$$du^i = 2u_3^i - u_i^i > 0, \quad i = 1, 2$$

▷ Pareto improvement (both better off)!

- Forced transfer by government t ($t < v^1, t < v^2$):

$$\frac{dc^i}{dt} = \frac{d[I^i - v^i]}{dt} = -1 - \frac{dv^i}{dt}$$

$$\frac{dc^3}{dt} = 2 + \frac{dv^1}{dt} + \frac{dv^2}{dt}$$

- New transfer equilibrium:

$$\frac{dv^1}{dt} = \frac{dv^2}{dt} = -1$$

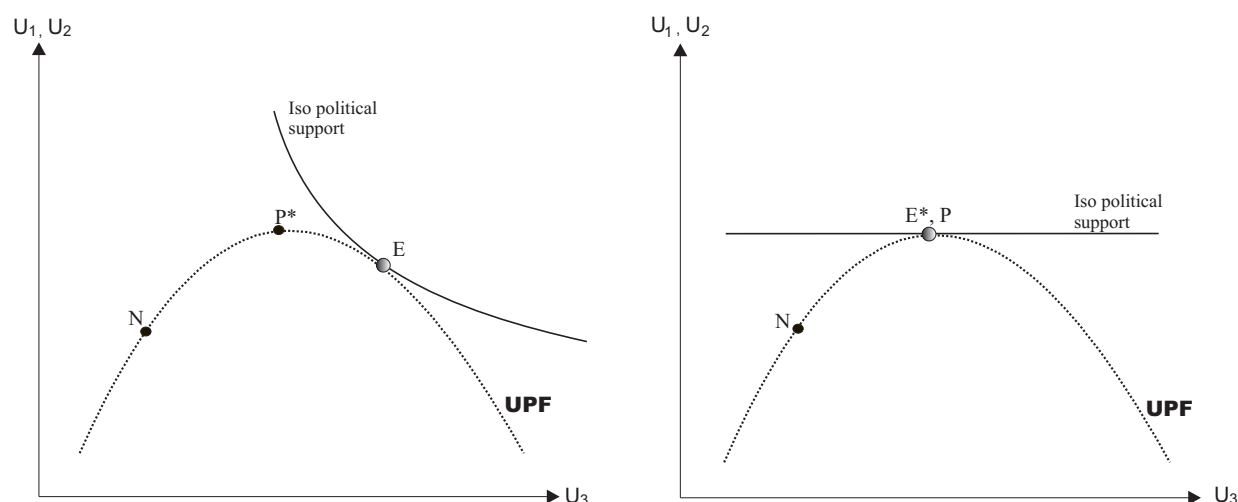
▷ Complete crowding-out, full offsetting

5.2. Political Support Model: Roberts [JPE 1984]

- Congress goal:

$$\max P(u_1, u_2, u_3)$$

- Stylized facts



5.3. Data Testing: Andreoni-Payne [JPuE 2011/v95]

- Two possible effects of government grants to charity:
 - Classical crowding-out: donors reduce contributions
 - Fundraising crowding-out: charities reduce fundraising efforts

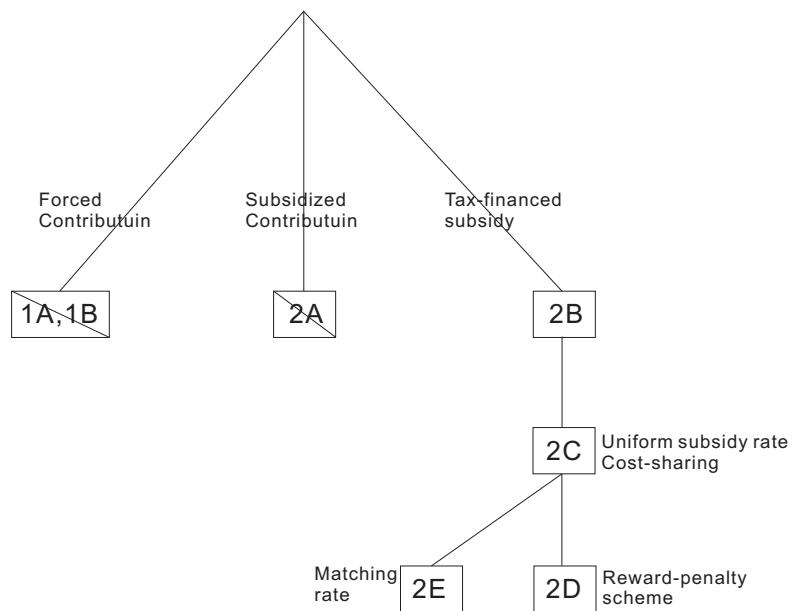
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- Panel data of tax returns from 8,000 charities:
 - Classical crowding-out: 30% to slight crowding-in
 - Fundraising crowding-out: 70–100%
 - Policy implication: requiring charities to match govt grants

6. Govt Financing/Subsidizing of PG Contribution

Notation:

- Consumer:
 - Wealth: m_i
 - Private consumption: c_i
 - Individual PG contribution: g_i
- Pre-tax Nash equilibrium:

$$(g_1^0, \dots, g_n^0)$$



6.1. Forced Contributions

1A Lump-sum Head Tax: Warr [JPuE 1982]

- Govt policy: individual head tax τ_i
- Total PG:

$$G = \sum_i g_i + \sum_i \tau_i$$

- Consumer budget:

$$c_i + g_i = m_i - \tau_i$$

- Full neutrality: unique Nash (if $\tau_i \leq g_i^0$)

$$c_i^* = c_i^0$$

$$g_i^* = g_i^0 - \tau_i$$

$$G^* = G^0$$

1B Labor Taxation: Bernheim [AER 1986]

- Consumer income: $m_i(l_i)$, $l_i \equiv$ individual labor supply
- Labor tax:

$$t_i(l), \quad l \equiv (l_1, \dots, l_n)$$

- Total PG:

$$G = \sum_i g_i + \sum_i t_i(l)$$

- Consumer budget:

$$c_i + g_i = m_i(l_i) - t_i(l)$$

- Pre-tax Nash equilibrium:

$$(g_1^0, \dots, g_n^0, l_1^0, \dots, l_n^0)$$

$$G^0 = \sum_i g_i^0$$

- Full neutrality:

$$c_i^* = c_i^0$$

$$l_i^* = l_i^0$$

$$g_i^* = g_i^0 - t_i(l^0)$$

$$G^* = G^0$$

▷ Unique Nash: for any $t_i(l)$ with

$$t_i(l^0) \leq g_i^0$$

6.2. Subsidized Contributions

2A Individual Subsidy Rate and Head Tax: Andreoni [JPuE 1988, Sec 3.1]

- Govt policy:
 - Subsidy rate: $\beta_i \in [0, 1]$
 - Fixed head tax: τ_i
- Govt PG addition: residual revenues

$$T = \sum_i \tau_i - \sum_i \beta_i g_i$$

- ▷ Government budget balanced
- ▷ Total PG:

$$G = \sum_i g_i + T = \sum_i [1 - \beta_i] g_i + \sum_i \tau_i$$

- Consumer budget:

$$c_i + [1 - \beta_i] g_i = m_i - \tau_i$$

! For \$1 contribution, consumer can still only raise G by 1

$$\text{MRS}_i^{G,x} = 1$$

- Full neutrality:¹ unique Nash for any $\{\beta_i, \tau_i\}$ with $\tau_i \leq g_i^0$

$$c_i^* = c_i^0, \quad g_i^* = \frac{g_i^0 - \tau_i}{1 - \beta_i}$$

$$G^* = \sum_i g_i^* + T = G^0$$

¹Note that the relevant part of i 's budget remains same if all other consumers act according to the offsetting rule. PG attributable to i is $[1 - \beta_i]g_i^* + \tau_i = g_i^0$.

2B Individual Subsidy Rate and Head Tax: A general framework

- *Tax-financed subsidy:*

- Individual subsidy rate:

$$\beta_i$$

- Linked head tax:

$$\tau_i$$

- Government budget balanced in equilibrium:

$$\{\beta_i, \tau_i\}, \quad \forall i$$

such that:

$$\sum_i \beta_i g_i = \sum_i \tau_i$$

▷ Total PG:

$$G = \sum_i g_i$$

- New equilibrium:

- Budget illusion: “myopic” consumers

- ▷ Consumers take (β_i, τ_i) as fixed/given

- Consumer budget:

$$c_i + [1 - \beta_i]g_i = m_i - \tau_i$$

- ▷ Consumer choice:

$$\text{MRS}_i = 1 - \beta_i$$

- For Samuelson efficiency:

$$\sum_i \text{MRS}_i = n - \sum_i \beta_i = 1$$

Hence:

$$\sum_i \beta_i = n - 1$$

[E] Equal cost sharing:

$$\beta_i = \frac{n-1}{n}, \quad 1 - \beta_i = \frac{1}{n} \quad \square$$

[N] For \$1 contribution, consumer can raise G by

$$\frac{1}{1 - \beta_i}$$

2C Uniform Subsidy and Cost Sharing: Andreoni-Bergstrom [PC 1996]

- Govt policy: *tax-financed subsidy*

– Uniform subsidy rate:

$$\beta \in [0, 1]$$

– Individual cost share:

$$s_i$$

– Individual tax total:

$$s_i \beta G$$

- Govt budget balanced:

$$\sum_i s_i = 1$$

- Total PG:

$$G = \sum_i g_i$$

- Consumer budget:

$$c_i + [1 - \beta]g_i = m_i - s_i \beta G$$

▷

$$c_i + (1 - \beta[1 - s_i])g_i = m_i - s_i \beta G_{-i}$$

N Special case of 2B:

$$\beta_i = \beta[1 - s_i], \quad \tau_i = s_i \beta G_{-i}$$

- Nash equilibrium: individual interior foc

$$\text{MRS}_i = 1 - \beta[1 - s_i]$$

- For Samuelson efficiency ($\sum \text{MRS} = 1$):

$$\sum_i \text{MRS}_i = n - n\beta + \beta = 1$$

▷ Just let

$$\beta = 1$$

Thm1 Unique Nash for any $\beta < 1$ and $\{s_i\}$

Thm2 G rises with β .

▷ \exists unique efficient G with $\beta = 1$, but $\{g_i\}$ not unique.

2D Reward/Penalty Scheme: Falkinger [JPuE 1996]

- Govt subsidy:

$$r_i = \beta[g_i - \bar{g}_i] \gtrless 0$$

where

$$\bar{g}_i \equiv \frac{G_{-i}}{n-1}$$

▷ Reward for contribution above average, penalty if below average.

- Govt budget always balanced, since $\sum_i r_i = 0$

- Total PG:

$$G = \sum_i g_i$$

- Consumer budget:

$$c_i + g_i = m_i + r_i$$

▷

$$c_i + [1 - \beta]g_i = m_i - \beta\bar{g}_i = m_i - \frac{\beta G_{-i}}{n-1}$$

N Special case of 2C: with $\tilde{\beta}$ and $s_i = 1/n$

$$\begin{aligned} c_i + g_i &= m_i + \beta[g_i - \frac{G_{-i}}{n-1}] \\ &= m_i + \frac{n-1}{n}\tilde{\beta}[g_i - \frac{G_{-i}}{n-1}], \quad \tilde{\beta} \equiv \frac{n}{n-1}\beta \\ &= m_i + \tilde{\beta}g_i - \frac{\tilde{\beta}}{n}G \end{aligned}$$

- Individual interior foc:

$$\text{MRS}_i = 1 - \beta$$

▷

$$\sum_i \text{MRS}_i = n - n\beta$$

▷ For efficiency: let

$$\beta = \frac{n-1}{n}$$

Thm1 If $0 < \beta < \frac{n-1}{n}$, then \exists unique Nash, and G^* increases with β .

Thm2 If $\beta = \frac{n-1}{n}$, then equilibrium not unique.

- Efficiency can be approximated with

$$\beta \rightarrow \frac{n-1}{n}$$

- Intuition: incentives for simultaneous cooperation in PD game.

2E Individual Subsidy and Tax Rates: *tax-financed subsidy*

- Govt policy:

– Individual subsidy rate:

$$\beta_i$$

– Linked matching rate:

$$t_i$$

- Consumer budget:

$$c_i + [1 - \beta_i]g_i = m_i - t_i G_{-i}$$

N Special case of 2B

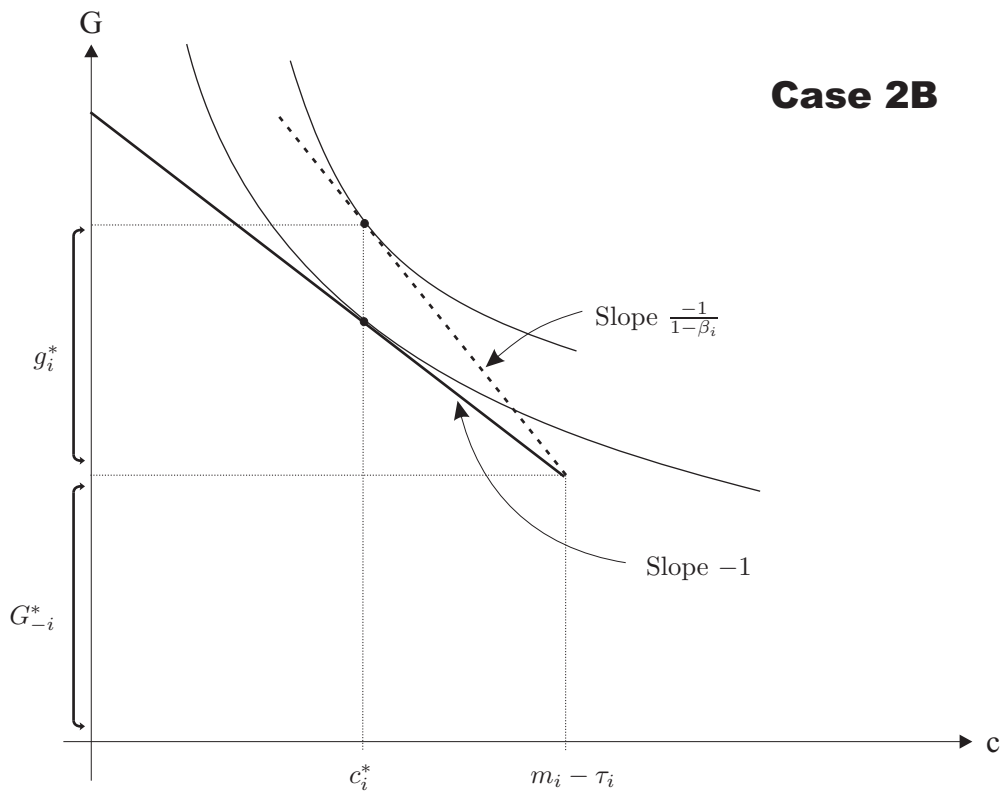
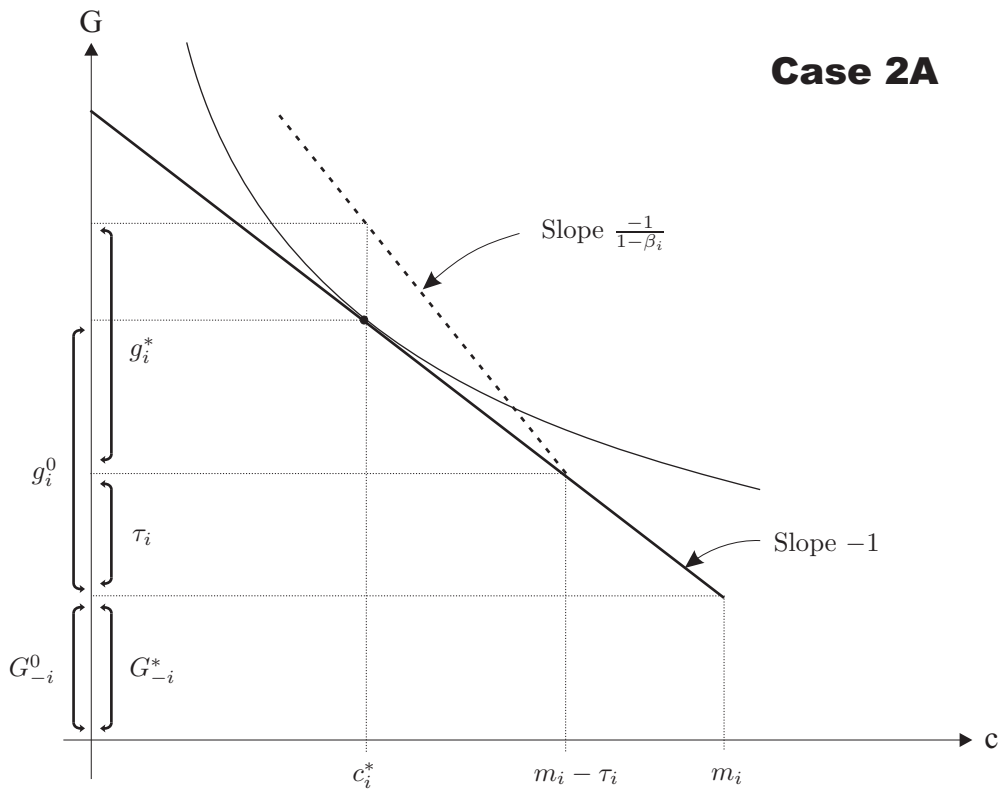
E Guttman [AER 1978]

- Govt budget balance in equilibrium:

$$\sum_i \beta_i g_i = \sum_i t_i G_{-i}$$

▷

$$\sum_i [\beta_i + t_i] g_i = \left[\sum_i t_i \right] G$$



7. The Rotten Kid Theorem (RKT)

7.1. Becker [1974, 1981]

- Household:

- Head: wealth I_H , consumption x_H

- Kids i ($= 1, \dots, n$): income I_i , consumption x_i

- Head's goal: max HH welfare with *non-negative* transfers

$$\max_{\{t_i\}} U_H(x_1, \dots, x_n, x_H) \quad \text{s.t.} \quad \begin{cases} x_H + \sum_{i=1}^n t_i = I_H \\ x_i = t_i + I_i \\ t_i \geq 0, \forall i \end{cases}$$

- Assuming *interior* solution ($t_i > 0$):

$$\max_{\{x_i\}} U_H(x_1, \dots, x_n, x_H) \quad \text{s.t.} \quad x_H + \sum x_i = I_H + \sum I_i$$

▷ *Social income*:

$$I_H + \sum_i I_i$$

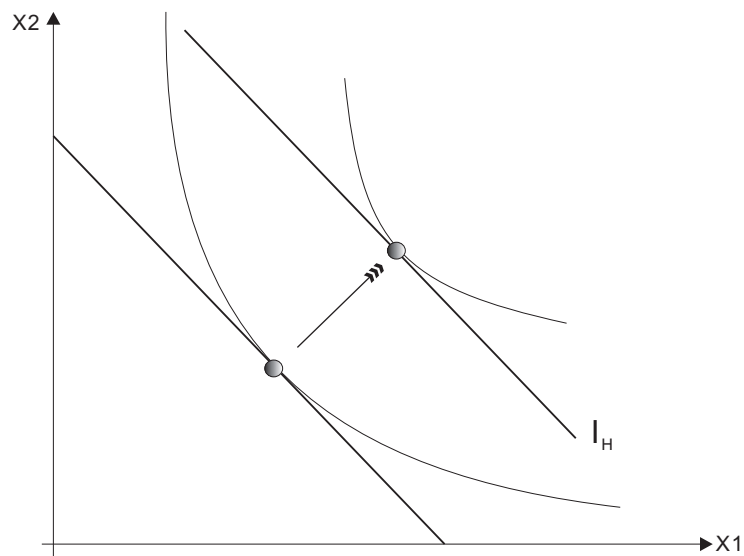
- HH interior foc:

$$\frac{\partial U_H}{\partial x_i} = \frac{\partial U_H}{\partial x_j}; \quad \forall i, j$$

- Kids will be *well-behaved*:

- Will maximize HH wealth ($I_H + \sum I_i$) for higher own x_i
- Will not steal from parents/siblings

□ E Concubines, party members



7.2. Bergstrom [1989]

Counter Examples

- Example #1: Lazy rotten kids

Kids: work y_i , leisure $[1 - y_i]$

$$U_i = x_i[1 - y_i]$$

Head:

$$\begin{aligned} \max_{x_1, x_2} U(x_1, x_2 | y_1, y_2) &\equiv \sqrt{U_1} + \sqrt{U_2} = \sqrt{x_1[1 - y_1]} + \sqrt{x_2[1 - y_2]} \\ \text{s.t. } x_1 + x_2 &= I_0 + w[y_1 + y_2] \end{aligned}$$

foc:

$$\frac{x_1}{x_2} = \frac{1 - y_1}{1 - y_2}$$

▷ No kid will want to work!

- Example #2: Night light (Becker [1974, p.1078])

Husband: income I , night reading hours y

$$u_h = x_h[y + 1]$$

Wife:

$$u_w = x_w \cdot e^{-y}$$

Husband's goal:

$$\max U = u_h \cdot u_w^a \quad \text{s.t.} \quad x_h + x_w = I$$

hence:

$$x_h^* = \frac{I}{1+a}, \quad x_w^* = \frac{aI}{1+a}, \quad y^* = \frac{1}{a} - 1$$

▷ Wife should cut the wire!

- Example #3: Prodigal son (Lindbeck-Weibull [JPE 1988])

2-period: $t = 1, 2$

Kid:

$$U_k = C_k^1 \cdot C_k^2$$

Head: max HH utility

$$U_H = [C_p^1 C_p^2] \cdot [U_k]^\alpha = [C_p^1 C_p^2] [C_k^1 C_k^2]^\alpha$$

FOC:

$$\frac{C_p^2}{C_k^2} = \frac{1}{\alpha}$$

▷ Kid consumes too much in period 1.²

²This is the Samaritan's Dilemma.

The 2-stage game Γ Rotten Kids Play

- n kids: individual action a_i ($\in A_i$)

$$a \equiv (a_1, a_2, \dots, a_n)$$

- Consumption of kid i :

$$m_i(a)$$

- HH budget:

$$M(a) \equiv \sum_i m_i(a)$$

- Kid utility:

$$u_i(a, m_i)$$

- Two stages:

- **S1** Kids choose a_i first
- **S2** Head makes transfer

$$\max_{\{m_i\}} U(u_1, u_2, \dots, u_n \mid a)$$

Analysis

- RKT holds if SPE of the game Γ accords with head's goal:

$$\max_{\{a_i, m_i\}} U(u_1(a, m_1), \dots, u_n(a, m_n))$$

▷ Kids are “well-behaved”

- The iff condition:

- Gorman utility function:

$$u_i(m_i, a) = A(a)m_i + B_i(a), \quad \forall i$$

- Utility possibility set (UPS): $UP(a)$ is a *simplex*

$$\sum_i u_i = A(a)M(a) + \sum_i B_i(a) = K(a)$$

▷ Transferable utility

- Examples revisited:

- Lazy rotten kids:

$$U_i(m_i, y) = m_i + B_i(y_i)$$

▷

$$UP(y_1, \dots, y_n) = \sum U_i = M(y_1, y_2, \dots, y_n) + \sum_i B_i(y_i)$$

- Night light:

$$u_h(m_h, y) = A(y)m_h + B_h(y)$$

$$u_w(m_w, y) = A(y)m_w + B_w(y)$$

▷

$$UP(y) = u_h + u_w = A(y)I + [B_h(y) + B_w(y)]$$

– Prodigal son:

$$U_i(C_i^1, C_i^2) = C_i^2 + B_i(C_i^1), \quad \forall i$$

▷

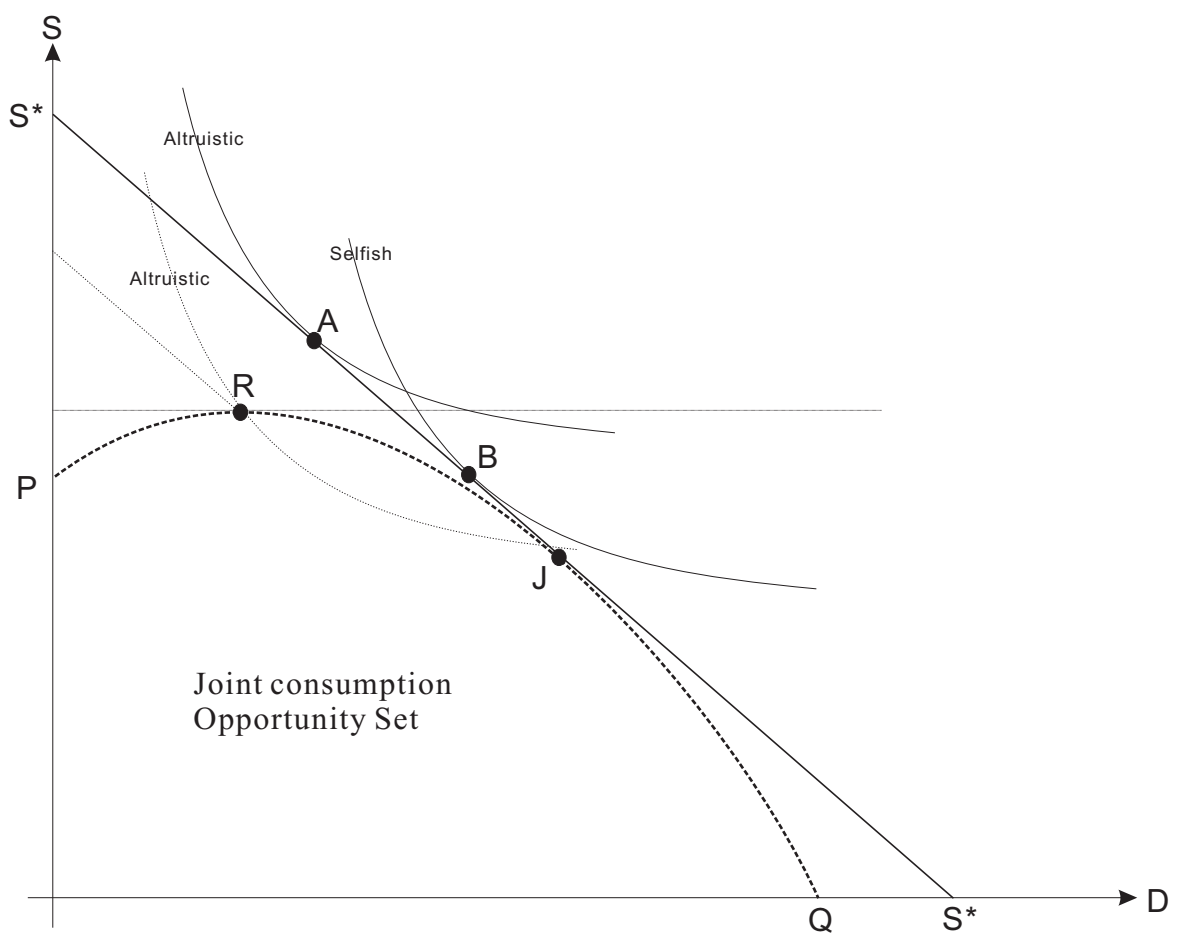
$$\text{UPF} = \sum_i C_i^2 + \sum_i B_i(C_i^1) = \left[I - \sum_i C_i^1 \right] + \sum_i B_i(C_i^1)$$

7.3. Bruce-Waldman [QJE 1990]

- 2-stage consumption-saving model
- Timing of transfer
- Samaritan's Dilemma: Buchanan 1975

7.4. Implication#1: Altruism in Evolution

- Biology: group rationality, genetic fitness, kin selection
- Becker: individual rationality³
- Hirshleifer [JEL 1977]



- Khalil [JEBO 2004, 53:89–92]

– Prosocial preferences:

³“Altruism, Egoism, and Genetic Fitness: Economics and Sociology,” *JEL*, G. Becker, 1976, 14(3):817–26.

- * Cooperation: for self-image⁴
- * Altruism: for other-interest⁵
- A group may not be better off with more altruists
- [E] Costly mutual help
- [E] Private taste/benefit info

7.5. Implication#2: United Way Donations

- UW will offset individual donation [Bilodeau, JPuE 1992]
 - ▷ You might as well give money to UW
- Contribution earmarking [Bilodeau-Slivinski, JPuE 1997]
 - ▷ Donor designations
 - ▷ Not to discourage individual donations.

⁴不好意思不做: 給小費, 讓座, 婚宴送禮, 履行交易承諾。

⁵甘心樂意: 給乞丐錢, 匿名捐獻, 服務志工。

8. Government by Jury

8.1. The Logic

- Voting paradox: why do people vote?

$$\text{Voting Costs} > \text{Expected Voting Benefits}$$

▷ *D*-term solution:

$$\text{Voting Costs} < \text{Expected Voting Benefits} + D$$

- Votes not intelligent?
 - ▷ PG argument
- Similar PG examples: providing public services
 - Congressman
 - Court judge
- Government by Jury: Bergstrom-Varian [1987]
 - Small congress
 - Random selection by computer
 - * Statistical (demographic) advantage
 - * Minimize social costs: election, gathering
 - * Argument against democracy v. GJ

8.2. Base Model: Constant Marginal Cost, Equal Benefit

- Population: m

Congress size: n

- Congressman:

– Effort:

$$e_i$$

– Constant MC:

$$c$$

– Total congress efforts:

$$E = \sum e_i$$

- Citizen benefits:

$$B(E); \quad B' > 0, \quad B'' < 0$$

- Optimality:

$$\max_E \pi \equiv mB(E) - cE$$

foc:

$$mB'(E^*) = c$$

8.3. Nash Equilibrium

- Individual foc:

$$B'(\hat{E}) = c$$

- Any combination of e_i for \hat{E} will do
- Symmetric solution:

$$\hat{e} = \frac{\hat{E}}{n}$$

$$B'(n \cdot \hat{e}) = c$$

- Total Nash effort is independent of n :⁶

$$\hat{e}(n) = \frac{\hat{e}(1)}{n}$$

- Severe free-riding
- To min TSC, should have few congressmen, picked by lottery
 - ▷ Lack of diversity

⁶Because

$$B'(n \cdot \hat{e}(n)) = B'(n \cdot \frac{\hat{e}(1)}{n}) = B'(\hat{e}(1)) = c$$

8.4. Modification#1: Different Benefit $B_i(E)$

- Nash: only one member k makes sole effort \hat{E} :

$$B'_k(\hat{E}) = c$$

$$B'_i(\hat{E}) \leq c, \forall i \neq k$$

- Free-riding more severe
 - Only the highest B_k individual makes sole effort $e_k = \hat{E}$
 - All others make no efforts $e_i = 0$

8.5. Modification#2: Increasing Marginal Costs

- Effort cost:

$$C(e_i), C' > 0, C'' > 0$$

- Symmetric Nash:

$$B'(n \cdot e(n)) = C'(e(n))$$

- Congress size effect:

$$\frac{de}{dn} = \frac{-eB''}{nB'' - C''} = \frac{-e}{n - \frac{C''}{B''}} < 0$$

$$\frac{dE}{dn} = \frac{d[n \cdot e(n)]}{dn} = \frac{e}{1 - n \cdot \frac{B''}{C''}} > 0$$

- Welfare effect:

$$\pi(n) = m \cdot B(n \cdot e(n)) - n \cdot C(e(n))$$

$$\left. \frac{d\pi(n)}{dn} \right|_{(m,n)} = m \cdot B'(E)E'(n) - [C(e) + nC'(e)e'(n)] \gtrless 0$$

▷ Positive if m is large enough, justifying large congress