

Binary System --Basics

The AND operation

- $0 \text{ AND } 0 = 0$
- $0 \text{ AND } 1 = 0$
- $1 \text{ AND } 0 = 0$
- $1 \text{ AND } 1 = 1$

The OR operation

- $0 \text{ OR } 0 = 0$
- $0 \text{ OR } 1 = 1$
- $1 \text{ OR } 0 = 1$
- $1 \text{ OR } 1 = 1$

The XOR operation

- $0 \text{ XOR } 0 = 0$
- $0 \text{ XOR } 1 = 1$
- $1 \text{ XOR } 0 = 1$
- $1 \text{ XOR } 1 = 0$

The NOT operation

- $\text{NOT } 0 = 1$
- $\text{NOT } 1 = 0$

$\begin{array}{r} 0 \\ + 0 \\ \hline 00 \end{array}$
 $\begin{array}{r} 0 \\ + 1 \\ \hline 01 \end{array}$
 $\begin{array}{r} 1 \\ + 0 \\ \hline 01 \end{array}$
 $\begin{array}{r} 1 \\ + 1 \\ \hline 10 \end{array}$

$\begin{array}{r} A \\ + B \\ \hline \end{array}$

Summary

\swarrow A and B \searrow A XOR B \downarrow NOT

A	B	AND	OR	XOR	NAND	NOR
0	0	0	0	0	1	1
0	1	0	1	1	1	0
1	0	0	1	1	1	0
1	1	1	1	0	0	0

The hexadecimal coding system

$$a_3 a_2 a_1 a_0 = a_3 \times 2^3 + a_2 \times 2^2 + a_1 \times 2^1 + a_0$$

- 0000 0 1000 8
- 0001 1 1001 9
- 0010 2 1010 10 **A**
- 0011 3 1011 11 **B**
- 0100 4 1100 12 **C**
- 0101 5 1101 13 **D**
- 0110 6 1110 14 **E**
- 0111 7 1111 15 **F**

Figure 1.7: The organization of a byte-size memory cell



Figure 1.14: The base ten and binary systems

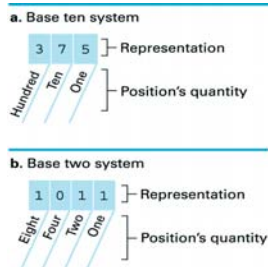
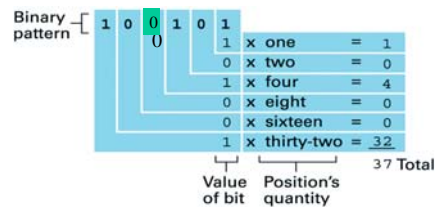


Figure 1.15: Decoding the binary representation 100101



An algorithm for finding the binary representation of a positive integer

- Step 1: Divide the value by two and record the remainder
- Step 2: As long as the quotient obtained is not zero, continue to divide the newest quotient by two and record the remainder.
- Step 3: Now that a quotient of zero has been obtained, the binary representation of the original value consists of the remainders listed from right to left in the order they were recorded.

Figure 1.17: Applying the algorithm in Figure 1.15 to obtain the binary representation of thirteen

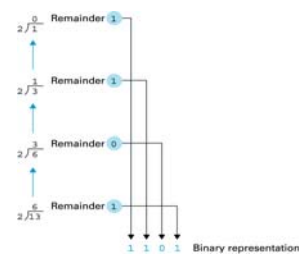


Figure 1.19: The binary addition facts

$$\begin{array}{r} 6 \\ + 8 \\ \hline 14 \end{array}$$

$$\begin{array}{r} 0 \quad 1 \quad 0 \quad 1 \\ +0 \quad +0 \quad +1 \quad +1 \\ \hline 0 \quad 1 \quad 1 \quad 10 \end{array}$$

Complement

- $67 - 55 = 67 - (100 - 45)$
- $= 67 + 45 - 100$
- $= 12$

Figure 1.21: Two's complement notation systems

a. Using patterns of length three		b. Using patterns of length four	
Bit pattern	Value represented	Bit pattern	Value represented
011	3	0111	7
010	2	0110	6
001	1	0101	5
000	0	0100	4
111	-1	0011	3
110	-2	0010	2
101	-3	0001	1
100	-4	0000	0
		1111	-1
		1110	-2
		1101	-3
		1100	-4
		1011	-5
		1010	-6
		1001	-7
		1000	-8

Figure 1.22: Coding the value -6 in two's complement notation using four bits

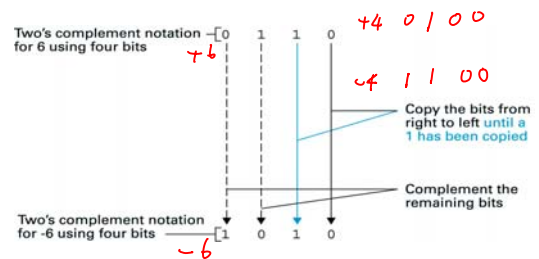


Figure 1.23: Addition problems converted to two's complement notation

Problem in base ten	Problem in two's complement	Answer in base ten
$\begin{array}{r} 3 \\ + 2 \end{array}$	$\begin{array}{r} 0011 \\ + 0010 \\ \hline 0101 \end{array}$	5
$\begin{array}{r} -3 \\ + -2 \end{array}$	$\begin{array}{r} 1101 \\ + 1110 \\ \hline 1011 \end{array}$	-5
$\begin{array}{r} 7 \\ + -5 \end{array}$	$\begin{array}{r} 0111 \\ + 1011 \\ \hline 0010 \end{array}$	2

Handwritten notes: $6 \ 0110$, $+7 \ 0111$, $(-3) \ 1101$

Overflow? E.g. $6+7=-3$ or $-6-8=(-6)+(-8)=+2$
Machine can make mistakes. It is treated with a special procedure.

Figure 1.24: An excess eight conversion table

Bit pattern	Value represented
1111	7
1110	6
1101	5
1100	4
1011	3
1010	2
1001	1
1000	0
0111	-1
0110	-2
0101	-3
0100	-4
0011	-5
0010	-6
0001	-7
0000	-8

5-tip

Figure 1.25: An excess notation system using bit patterns of length three

Bit pattern	Value represented
111	3
110	2
101	1
100	0
011	-1
010	-2
001	-3
000	-4

Figure 1.20: Decoding the binary representation 101.101

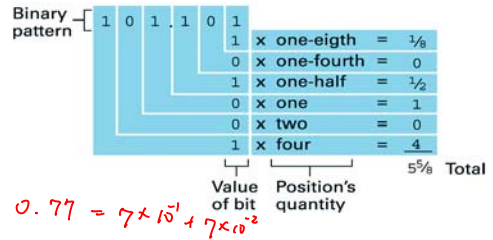


Figure 1.26: Floating-point notation components

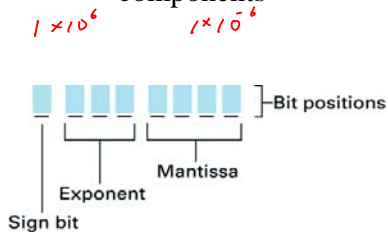


Figure 1.27: Coding the value 2 5/8

