

Chapter 1 Mathematical Preliminaries and Error Analysis • Ex. 證明 $0.\overline{9} = 1$ Proof: Let $x = 0.\overline{9}$ $10x = 9.\overline{9}$ 9x = 9 $\Rightarrow x = 1$

- This book examines problems that can be solved by methods of a approximation, which we call numerical methods. For example, numerical differential.
- Ex. $f(x) = x^5 + 5x^4 + 3x^3 2x^2 + x 1$, 求 $f(\pi)$. Sol: ((((x+5)x+3)x-2)x+1-1)(因 trumcate or round off 之次數少, 所以誤差小)
- Continuity and Differentiability (P.2~P.3)

1.2 Review of Calculus

• <u>Mean Value Thorem</u>

If $f \in C[a, b]$ and f is differentiable on (a, b), then a number c in (a, b) exists such that (see P.4 Figure 1.3)

 $f'(c) = \frac{f(b) - f(a)}{b - a}$

• <u>Extreme Value Thorem</u>

If $f \in C[a, b]$, then c_1 and c_2 in [a, b] exist with $f(c_1) \leq f(x) \leq f(c_2)$ for all x in [a, b]. If, in addition f is differentiable on (a, b), then the number c_1 and c_2 occur either at endpoints of [a, b] or where f' is zero.

(極値不是在邊界, 就是在 f' = 0)

• Mean Value Thorem for Integrals and Intermediate Value Theorem

If $f \in C[a, b]$, g is integrable on [a, b] and g(x) does not change sign on [a, b], then there exists a number c in (a, b)with $\int_a^b f(x)g(x)dx = f(c)\int_a^b g(x)dx$ (可找到一個等面積之長方形)

• Taylor's Thorem

Suppose $f \in C^n[a, b]$ and f^{n+1} exists on [a, b]. Let x_0 be a number in [a, b]. For every x in [a, b], there exists a number $\xi(x)$ between x_0 and x with

 $f(x) = P_n(x) + R_n(x)$

where

$$P_n(x) = f(x_0) + f'(x_0)(x - x_0) + \frac{f''(x_0)}{2!}(x - x_0)^2 + \dots + \frac{f^n(x_0)}{n!}(x - x_0)^n = \sum_{k=0}^n \frac{f^k(x_0)}{k!}(x - x_0)^k$$

and

$$R_n(x) = \frac{f^{n+1}(\xi(x))}{(n+1)!} (x - x_0)^{n+1}$$

(用 x₀ 這個點上的各階動差來找出整個函數)
(P.9 Figure 1.7)

• $f(x,y) = f(x_0,y_0) + \frac{\partial f}{\partial x}(x-x_0) + \frac{\partial f}{\partial y}(y-y_0) + \frac{1}{2!}\frac{\partial^2 f}{\partial x^2}(x-x_0)^2 + \frac{1}{2!}\frac{\partial^2 f}{\partial y^2}(y-y_0)^2 + \frac{\partial^2 f}{\partial x \partial y}(x-x_0)(y-y_0)$ (二維) (Itô's Lemma)

1.3 Round-off Error and Computer Arithmetic

• To save storage and provide a unique representation for each floating-point number, a normalization is imposed. Using this system gines a floating-point number of the form

$$(-1)^s \times 2^{c-1023} \times (1+f)$$

Consider for example, the machine number

(s c f) = (sign characteristic mantissa)

The leftmost bit is zero, which indicates that the number is positive. The next 11 bits, 10000000011, giving the characteristic, are equivalent to the decimal number

 $c = 1 \times 2^{10} + 0 \times 2^9 + \ldots + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 = 1024 + 2 + 1 = 1027$

The exponential part of the number is, therefore, $2^{1027-1023} = 2^4$. The final 52 bits specify that the mantissa is

$$f = 1 \times (\frac{1}{2})^1 + 1 \times (\frac{1}{2})^3 + 1 \times (\frac{1}{2})^4 + 1 \times (\frac{1}{2})^5 + 1 \times (\frac{1}{2})^8 + 1 \times (\frac{1}{2})^{12}$$

As a consequence, this machine number precisely represents the decimal number

$$(-1)^{s} \times 2^{c-1023} \times (1+f) = (-1)^{0} \times 2^{1027-1023} (1 + (\frac{1}{2} + \frac{1}{8} + \frac{1}{16} + \frac{1}{32} + \frac{1}{256} + \frac{1}{4096})) = 27.56640625$$

However, the next smaller machine number is

and the next larger machine number is

This means that our original machine number represents not only 27.56640625, but also half of the real numbers that are between 27.56640625 and its two nearest machinenumber neighbors. • The smallest normalized positive number that can be represented has all zeros except for the rightmost bit of 1 and os equivalent to

 $2^{-1023}(1+2^{-52}) \approx 10^{-308}$

and the largest has a leading 0 followed by all 1s and is equivalent to

$$2^{1024}(2-2^{-52}) \approx 10^{308}$$

Numbers occuring in calculations that have a magnitude less than $2^{-1023}(1+2^{-52})$ result in <u>underflow</u> and are generally set to zero . Numbers greater than $2^{1024}(2-2^{-52})$ result in <u>overflow</u> and typically cause the computation to halt.

• Rounding and Chopping

(The error that results from replacing with its floating-point form is called round-off error)

1.4 Errors in Scientific Computation

• (P.21)

Ex 1. $x^2 + 62.10x + 1 = 0$, whose roots are approximately $x_1 = -0.0160723$ and $x_2 = -62.08390$

In this equation, b^2 is much larger than 4ac, so the numerator in the calculation for x_1 involves the subtration of nearly equal numbers. Since

$$\sqrt{b^2 - 4ac} = \sqrt{(62.10)^2 - (4)(1)(1)} = 62.06 \ (62.06778552)$$

we have

 $fl(x_1) = \frac{-62.10+62.06}{2.000} = \frac{-0.04000}{2.000} = -0.02000$ (因 62.10 與 62.06 很接近, 使得 truncate 掉的 0.00778552 顯得很重要),

a poor approximation to $x_1 = -0.01611$ with the larger relative error

$$\frac{|-0.01611+0.002000|}{|-0.01611|} = 2.4 \times 10^{-1} \text{ (誤差很大)}$$

On the other hand, the calculation for x_2 involves the addition of nearly equal numbers -b and $-\sqrt{b^2 - 4ac}$. The presents no problem since

 $fl(x_2) = \frac{-62.10-62.06}{2.000} = \frac{-124.2}{2.000} = -62.10,$ has the small relative error $\frac{|-62.08+62.10|}{|-62.08|} = 3.2 \times 10^{-4}$ (誤差很小) 要獲得更精確的 x_1 逼近值,使用反有理化. $x_1 = \frac{-b+\sqrt{b^2-4ac}}{2a} = \frac{-2c}{b+\sqrt{b^2-4ac}}$ $\Rightarrow fl(x_1) = \frac{-2.000}{62.10+62.06} = -0.01610$ which has the small relative error 6.2×10^{-4} (反有理化後,誤差變小) 但同時, $fl(x_2)$ 若也做反有理化, 會從誤差小 \longrightarrow 誤差大.