The 2x2 Exchange Economy	
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(Lecture 7, Micro Theory I)	

Road Map for Chapter 3



- Pareto Efficiency
 - Cannot make one better off without hurting others
- Walrasian (Price-taking) Equilibrium
 - When Supply Meets Demand
 - Focus on Exchange Economy First
- 1st Welfare Theorem: Walrasian Equilibrium is Efficient (Adam Smith Theorem)
- 2nd Welfare Theorem: Any Efficient Allocation can be supported as a Walrasian Equilibrium

2x2 Exchange Economy

- 2 Commodities: Good 1 and 2
- 2 Consumers: Alex and Bev h = A, B
 - Endowment: $\omega^h = (\omega_1^h, \omega_2^h), \, \omega_i = \omega_i^A + \omega_i^B$
 - Consumption Set: $x^h = (x_1^h, x_2^h) \in \mathbb{R}^2_+$
 - Strictly Monotonic Utility Function:
- Edgeworth Box

 $U^{h}(x^{h}) = U^{h}(x_{1}^{h}, x_{2}^{h})$

• These consumers could be representative agents, or literally TWO people (bargaining)

Why do we care about this?



- The Walrasian (Price-taking) Equilibrium (W.E.) is (a candidate of) Adam Smith's "Invisible Hand"
 - Are real market rules like Walrasian auctioneers?
 - Is Price-taking the result of competition, or competition itself?
- Illustrate W.E. in more general cases
 - Hard to graph "N goods" as 2D
- Two-party Bargaining
 - This is what Edgeworth really had in mind

Why do we care about this?



- Consider the following situation: You company is trying to make a deal with another company
 - Your company has better technology, but lack funding
 - Other company has plenty of funding, but low-tech
- There are "gives" and "takes" for both sides
- Where would you end up making the deal?
 - Definitely not where "something is left on the table."
- What are the possible outcomes?
 - How did you get there?

Social Choice and Pareto Efficiency

Benthamite:

- Behind Veil of Ignorance U_2
- Assign Prob. 50-50

 $\max \frac{1}{2}U^A + \frac{1}{2}U^B$

- Rawlsian:
 - Extremely Risk Averse $\max \min\{U^A, U^B\}$
- Both are Pareto Efficient
 - But A is not

$A \circ U^A + U^B$



Pareto Efficiency



- A feasible allocation is Pareto efficient if
- there is no other feasible allocation that is
- strictly preferred by at least one consumer
- and is weakly preferred by all consumers.



Pareto Efficient Allocations

For
$$\omega = (\omega_1, \omega_2)$$
, consider

$$\max_{x^A, x^B} \left\{ U^A(x^A) | U^B(x^B) \ge U^B(\overline{x}^B), x^A + x^B \le \omega \right\}$$
Need $MRS^A(\hat{x}^A) = MRS^B(\overline{x}^B)$ (interior solution)

$$U^{B}(x^{B}) = U^{B}(\overline{x}^{B})$$

$$\overline{x}^{B}$$

$$\hat{x}^{A}$$

$$U^{A}(x^{A}) = U^{A}(\hat{x}^{A})$$

$$O^{A}$$

$$X_{1}$$

$$U^{A}(x^{A}) = U^{A}(\hat{x}^{A})$$

$$U^{A}(x^{A}) = U^{A}(\hat{x}^{A})$$

Walrasian Equilibrium (in 2x2 Exchange Economy)

- All Price-takers: Prices $p \ge 0$
- 2 Consumers: Alex and Bev h = A, B
 - Endowment: $\omega^h = (\omega_1^h, \omega_2^h), \, \omega_i = \omega_i^A + \omega_i^B$
 - Consumption Set: $x^h = (x_1^h, x_2^h) \in \mathbb{R}^2_+$

• Wealth:
$$W^h = p \cdot \omega^h$$

- Market Demand: $x(p) = \sum_{h} x^{h}(p, p \cdot \omega^{h})$
- Vector of Excess Demand: $e(p) = x(p) \omega$

• Vector of total Endowment: $\omega = \sum \omega^h$

h

Definition: Market Clearing Prices



- Let excess demand for commodity *j* be $e_j(p)$
- The market for commodity *j* clears if $e_j(p) \le 0$ and $p_j \cdot e_j(p) = 0$
- Why is this important?
- Walras Law
 - The last market clears if all other markets clear
- Market clearing defines Walrasian Equilibrium

Walras Law



- LNS implies consumer must spend all income
- If not, we have $p \cdot x^h$
- But then there exist δ -neighborhood $N(x^h, \delta)$
- In the budget set for sufficiently small $\delta > 0$
 - Contradicting LNS

$$\sum_{h} (p \cdot x^{h} - p \cdot \omega^{h}) = 0 = p \cdot \left(\sum_{h} (x^{h} - \omega^{h})\right)$$

 $= p \cdot (x - \omega) = p \cdot e(p) = p_1 e_1(p) + p_2 e_2(p) = 0$

• If one market clears, so must the other.

Definition: Walrasian Equilibrium



- The price vector $p \ge 0$ is a Walrasian Equilibrium price vector if all markets clear.
 - WE = price vector!!!
- EX: Excess supply of commodity 1...



Definition: Walrasian Equilibrium



- Lower price for commodity 1 if excess supply
- Until Markets Clear x_2 $O^B = (\omega_1, \omega_2)$ (ω_1^A, ω_2^A) x_1 Cannot raise Alex's utility without hurting Bev
 - Hence, we have...

First Welfare Theorem: WE \rightarrow PE



- If preferences satisfy LNS, then a Walrasian Equilibrium allocation (in an exchange economy) is Pareto efficient.
- Sketch of Proof:
- 1. Any weakly (strictly) preferred bundle must cost at least as much (strictly more) as WE
- 2. Markets clear
 - \rightarrow Pareto preferred allocation not feasible

First Welfare Theorem: WE \rightarrow PE



- 1. Since WE allocation \overline{x}^h maximizes utility, so $U^h(x^h) > U(\overline{x}^h) \Rightarrow p \cdot x^h > p \cdot \overline{x}^h$ Now need to show that $U^h(x^h) \ge U(\overline{x}^h) \Rightarrow p \cdot x^h \ge p \cdot \overline{x}^h$
- If not, we have $p \cdot x^h$
- But then there exist δ -neighborhood $N(x^h, \delta)$
- In the budget set for sufficiently small $\delta>0$
- Contradicts LNS that requires a point \tilde{x}^h such that $U^h(\tilde{x}^h) > U^h(x^h) \ge U(\overline{x}^h)$

First Welfare Theorem: WE \rightarrow PE



- 1. $U^h(x^h) > U(\overline{x}^h) \Rightarrow p \cdot x^h > p \cdot \overline{x}^h$ $U^h(x^h) \ge U(\overline{x}^h) \Rightarrow p \cdot x^h \ge p \cdot \overline{x}^h$ 2. True for Pareto preferred allocation (x^A, x^B) • Hence, $p \cdot x^h > p \cdot \overline{x}^h$ for at least one, and $p \cdot x^h > p \cdot \overline{x}^h$ for all others (by PEA) • Thus, $p \cdot \sum x^h > p \cdot \sum \overline{x}^h = p \cdot \sum \omega^h$ h_{-} hh• Since $p \ge 0$, at least one $j \rightarrow \sum x_i^h > \sum \omega_i^h$
 - Not feasible! $h = h^{20}$

Second Welfare Theorem: $PE \rightarrow WE$



- For a Pareto efficient allocation (\hat{x}^A, \hat{x}^B)
- Convex preferences imply convex regions
 - Separating hyperplane theorem generates prices



Second Welfare Theorem: $PE \rightarrow WE$



- If preferences are convex & strictly increasing, then any Pareto efficient allocation (of an exchange economy) can be supported by a price vector $p \ge 0$ (as a Walrasian Equilibrium).
- Sketch of Proof:
- 1. The at-least-as-good-as sets are convex
- 2. Supporting Hyperplane Theorem $\rightarrow p \ge 0$
- 3. Alex and Bev are both maximizing at $p \ge 0$



- Alex's "at-least-as-good-as" set X^A is convex
- Bev's "at-least-as-good-as" set X^B is convex
- $X^A \cap X^B$ has no interior since (\hat{x}^A, \hat{x}^B) is PEA



- By Supporting Hyperplane Theorem, exists *p*,
- Such that $p \cdot x^A > p \cdot \hat{x}^A \quad \forall x^A \in \text{int} X^A$ $p \cdot x^A$



- Claim $p \ge 0$, then $p \cdot x^A > p \cdot \hat{x}^A \quad \forall x^A \in \text{int} X^A$
- Implies Alex is maximizing. Similarly,
- Bev is maximizing since $x^A = \omega x^B \in \text{int} X^B$ $p \cdot x^B = p \cdot (\omega - x^A) > p \cdot (\omega - \hat{x}^A) = p \cdot \hat{x}^B$ ²⁵



Proposition 3.1-3 PEA with Homothetic Preferences

- 2x2 Exchange Economy
- Consumers have homothetic preferences
- At aggregate endowment, consumer A has a stronger preference for commodity 1.
- Consumption ratio: $\frac{x_2^A}{x_1^A} < \frac{x_2^B}{x_1^B}$
- And, as $U^A(x^A)$ rises, consumption ratio $\frac{x_2^A}{x_1^A}$ and MRS both rise.

Homothetic Preferences: Radial Parallel



- MRS same on each ray
- MRS increases as slope of the ray increase



Assumption: Intensity of Preferences



- Alex and Bev: convex, homothetic preferences
- Alex has stronger preferences for commodity 1







Proposition 3.1-3 PEA with Homothetic Preferences

- 2x2 Exchange Economy
- Consumers have homothetic preferences
- At aggregate endowment, consumer A has a stronger preference for commodity 1.
- Then at any interior PEA, $\frac{x_2^A}{x_1^A} < \frac{x_2^B}{x_1^B}$



• Moreover, as $U^A(x^A)$ rises, consumption ratio $\frac{x_2^A}{x_1^A}$ and MRS both rise.

Summary of 3.1



- Pareto Efficiency:
 - Cannot make one better off without hurting others
- Walrasian Equilibrium: market clearing prices
- Welfare Theorems:
 - First: Walrasian Equilibrium is Pareto Efficient
 - Second: Pareto Efficient allocations can be supported as Walrasian Equilibria (with transfer)
- Homework: Riley 3.1-1, 2, 4
- J/R 5.11, 5.12, 5.15, 5.17, 5.18