Games with Private Information 資訊不透明賽局

> Joseph Tao-yi Wang 2010/10/15

(Lecture 9, Micro Theory I-2)

Market Entry Game with Private Information















BNE when p > 1/2: (Strong Entrant *Enter*; Others Mix) $U^{1}(Enter; Weak) = \beta \cdot (-2) + (1 - \beta) \cdot 3$ $= 3-5\beta = U^1(Out;Weak) = 0$ 1-α **(0,6)** Fight (-2,4) if $\beta = \frac{3}{5}$ Out Enter 2 α **Share** (3,3) $1 - \beta$ Weak '-1,2) **Fight** Strong Enter 2 Share (3,3) $1 - \beta$ Out (0,6) Player1 will mix if Weak





BNE when p > 1/2: (Strong Entrant *Enter*; Others Mix) $U^2(Fight) = (1-q) \cdot 4 + q \cdot 2$ $= 4 - 2q = U^2(Share) = 3$ $1 - \alpha$ (0,6) Enter 2 Fight (-2,4) if $q = \frac{1}{2}$ α **Share** (3,3) $1 - \beta$ a =(1 - p)Weak $p \cdot \alpha + (1-p)$ Fight (-1,2) Strong Enter 2 Share (3,3) Out (0,6)Player 2 mixes if $\alpha = \frac{1-p}{r}$



Modified Market Entry Game: New Payoffs if Enter...











Separating Equilibrium: (Strong-*Enter*; Weak-*Out*) $U^2(Share) = 1 > U^2(Fight) = -6$ (0,6) Enter 2 Fight (-6,5) Out Share (-1,4) Weak Fight (-1,-6) Strong Enter 2 Share (4,1) Out (0,6)Player 2 will Share



(Strong-*Enter*; Weak-*Out*) is also a Sequential Equilibrium! $P(Weak|Enter) = \frac{p\epsilon_W}{(1-p)(1-\epsilon_S)+p\epsilon_W} \to 0$ as $\epsilon_W \to 0$ $1 - \epsilon_W$ Fight (-6,5) Enter Out 2 ϵ_W Share (-1,4) Weal Fight (-1,-6) Strong - p $1 - \epsilon_S$ Enter 2 Share (4,1) Out ϵ_S (0,6) $P(Weak|Enter) = \frac{P(Weak,Enter)}{P(Strong,Enter) + P(Weak,Enter)}$





Pooling Equilibrium: (Out, Out, Fight) $U^1(Out; Strong) = 0$ $> U^1(Enter; Strong) = -1$ (0,6) Fight (-6,5) Enter 2 Out Share (-1,4) Weak Fight (-1,-6) **Strong** -pEnter 2 Share (4,1) Out (0,6)Player1 stays Out if Strong





(Out, Out, Fight) is also a **Sequential Equilibrium!** (0,6) If $\epsilon_S = \frac{p}{1-p} \theta \epsilon_W$, P(Weak | Enter)1 - ϵ_W Out Enter 2 Fight (-6,5) = $\frac{p \epsilon_W}{(1-p) \epsilon_S + p \epsilon_W}$ **Share** (-1,4) $= \frac{1}{1+\theta}$ ϵ_W Weal Still Sequential!! Fight (-1,-6) Strong ϵ_S Enter 2 Share (4,1) (0,6) $P(Weak|Enter) = \frac{P(Weak,Enter)}{P(Strong,Enter) + P(Weak,Enter)}$

(*Out*, *Out*, *Fight*) is also a Sequential Equilibrium!



- (Out, Out, Fight) is not ruled out by THP, and hence, is also a Sequential Equilibrium...
- But why can't the Strong type say,
- "If I enter, I will be credibly signaling that I am *Strong*, since if I were weak and chose to *Enter*, my possible payoffs would be -1 or -5, smaller than 0 (equilibrium payoff if weak)."
- Seeing this, player 2's BR is *Share*
 - It is profitable for player 1 to *Enter* (& signal)...

Definition: Intuitive Criterion (Cho and Kreps)



- Consider \hat{a}_i , a strategy of player *i* that is not chosen in the Bayesian Nash equilibrium,
- Let $u_i(\hat{a}_i, t_i)$ be the payoff of player *i*'s if he chooses \hat{a}_i and is believed to be type $t_i \in T_i$
- Let $u_i^N(t_i)$ be this types' equilibrium payoff
- The BNE fails the Intuitive Criterion if, for some player *i* of type $\hat{t}_i \in T_i$, $u_i(\hat{a}_i, \hat{t}_i) > u_i^N(\hat{t}_i)$
- And for all other types in $t_i \in T_i$,

$$u_i(\hat{a_i}, t_i) < u_i^N(t_i)$$

Intuitive Criterion (Cho and Kreps)



- In the previous Example,
- (Out, Out, Fight) fails the Intuitive Criterion
 - "If I enter, I will be credibly signaling that I am Strong, since if I were weak and chose to Enter, my possible payoffs would be -1 or -5, smaller than 0 (equilibrium payoff if weak)."
- (*Out*, *Enter*, *Share*) satisfies the Intuitive Criterion
 - Such argument is not credible...

Continuous Types: An Auction Game



- One single item for sale
- *n* risk neutral bidders
- Valuation is continuously distributed on the unit interval with cdf $F(.) \sim [0,1]$
 - All this is common knowledge
- Bidder's type = Valuation (private information)
- Pure Strategy = Bid function $b = b_j(v_j)$

Sealed High-Bid Auction (aka First Price Auction)



• Each buyer submits one sealed bid

 $b_j \ge 0, j = 1, \cdots, n$

- Buyer who makes highest bid is the winner
 - If there is a tie, the winner is chosen randomly from the tying high bidders
- The winning bidder pays his bid and receivers the item

Sealed High-Bid Auction (aka First Price Auction)



- Bidder j, j=1, ..., n, knows own valuation v_j
 - Risk neutral, pay b, wins with probability p
- Payoff is $u_j(b, p, v_j) = p(v_j b)$
- Solve for Equilibrium Bidding Strategy $b_j(v_j)$
- For the special case of <u>2 bidders</u> of Independent Private Value (IPV)
- Assume valuation is uniform [0,1], cdf F(x) = x

BR to a Linear Strategy



- If buyer 2's bidding strategy is $b_2(v_2) = \alpha_2 v_2$
- Then the distribution of bids is uniform
 - Since valuation is uniform
- If buyer bids b, he wins with probability $p_2(b) = Pr[b_2 \le b] = \frac{b}{\alpha_2}, b \in [0, \alpha_2]$ • Buyer 1: $U_1(b) = p_2(b)(v_1 - b)$ $= \frac{b}{\alpha_2}(v_1 - b) = \frac{1}{\alpha_2}(bv_1 - b^2)$

Equilibrium of the Sealed High-Bid (aka First Price) Auction

• Solve the maximization problem:

$$\max_{b} U_1(b) = p_2(b)(v_1 - b) = \frac{1}{\alpha_2}(bv_1 - b^2)$$

• FOC:
$$v_1 - 2b = 0$$

• Maximum at
$$b_1(v_1) = \frac{1}{2}v_1$$

I.e. The BR to a linear strategy is a linear strategy

1

• By symmetry, the BNE is $b_j(v_j) = \frac{1}{2}v_j$

3 Bidder Case



- What if there are 3 bidders?
- Intuition is you would bid higher (competition)
- Assume bidder 3 bids $b_3(v_3) = \alpha_3 v_3$
- If buyer bids *b*, he wins with probability $p_3(b) = Pr[b_3 \le b] = \frac{b}{\alpha_3}, b \in [0, \alpha_3]$
- Buyer 1: $U_1(b) = p_2(b)p_3(b)(v_1 b)$ = $\frac{b}{\alpha_2} \frac{b}{\alpha_3}(v_1 - b) = \frac{1}{\alpha_2\alpha_3}(b^2v_1 - b^3)$

Equilibrium of the Sealed High-Bid (aka First Price) Auction

• Solve the maximization problem:

 $\max_{b} U_1(b) = p_2(b)p_3(b)(v_1 - b) = \frac{1}{\alpha_2 \alpha_3} (b^2 v_1 - b^3)$

• FOC:
$$2bv_1 - 3b^2 = 0$$

• Maximum at $b_1(v_1) = \frac{2}{3}v_1$

• I.e. The BR to a linear strategy is a linear strategy

• By symmetry, the BNE is $b_j(v_j) = \frac{2}{3}v_j$

Summary of 9.7



- Pooling Equilibrium vs. Separating Equilibrium
- Semi-Pooling Equilibrium (MSE)
- Intuitive Criteria
- Continuous Type Models: Auction Games
- HW 9.7: Riley 9.7-1~3