Games and Strategic Equilibrium

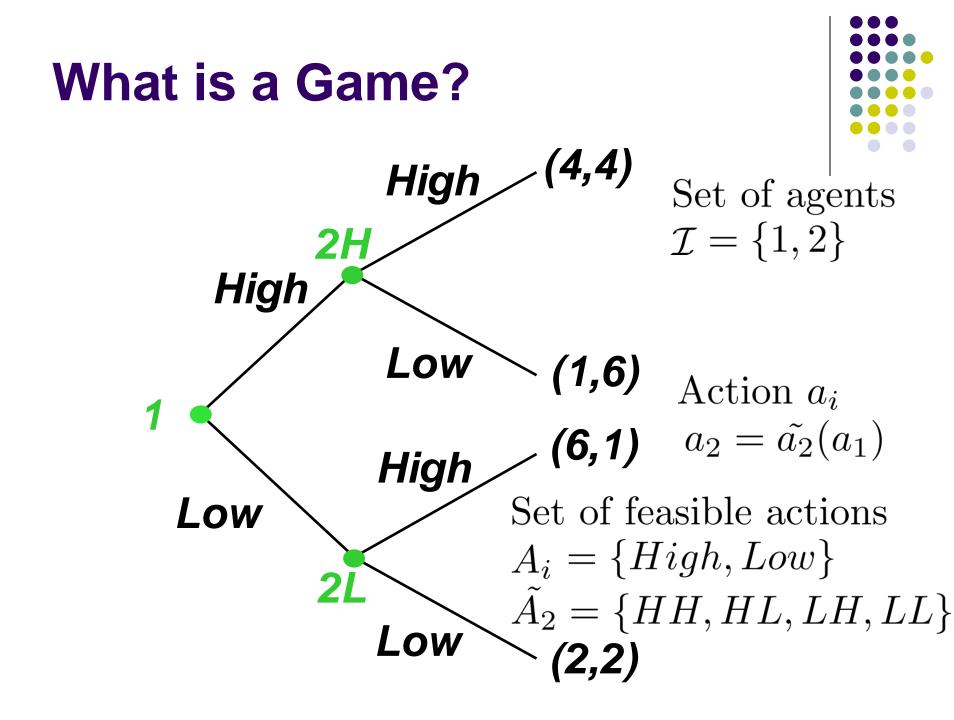
Joseph Tao-yi Wang 2010/9/17

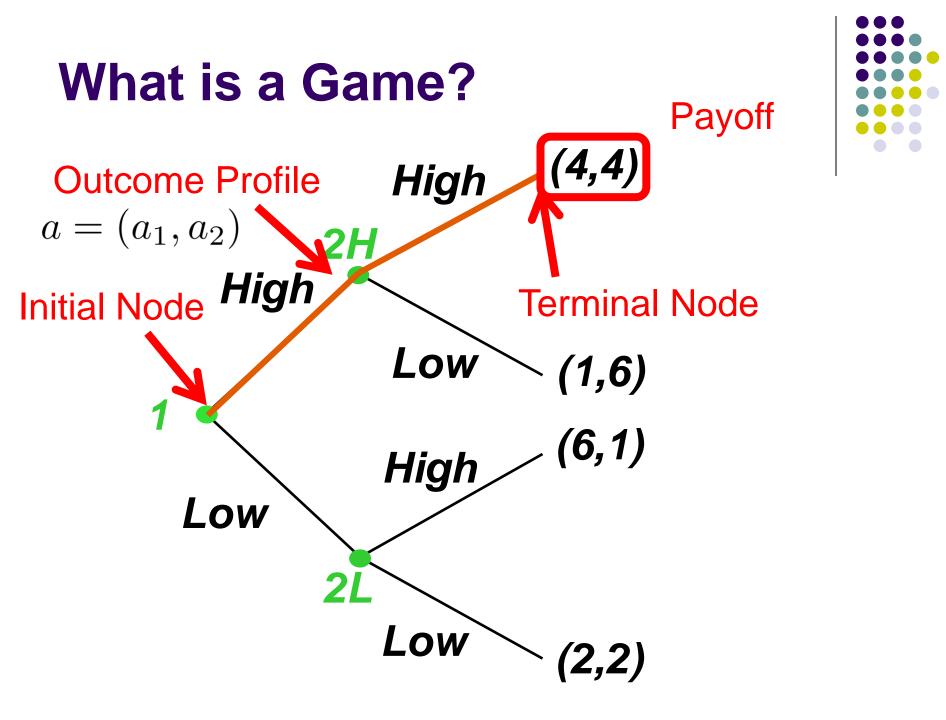
(Lecture 2, Micro Theory I-2)

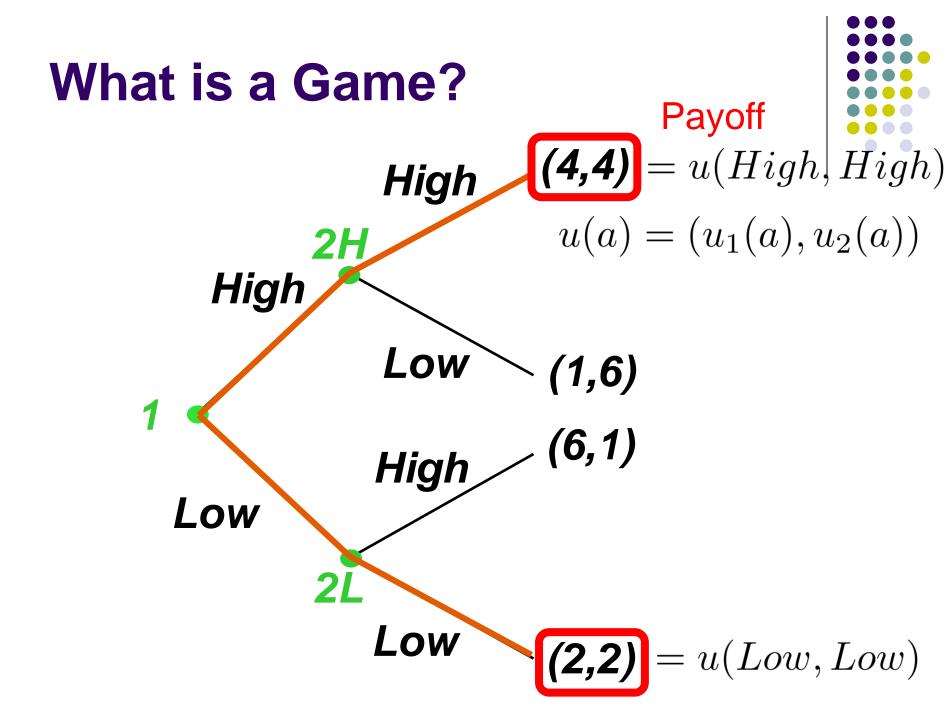
What is a Game?

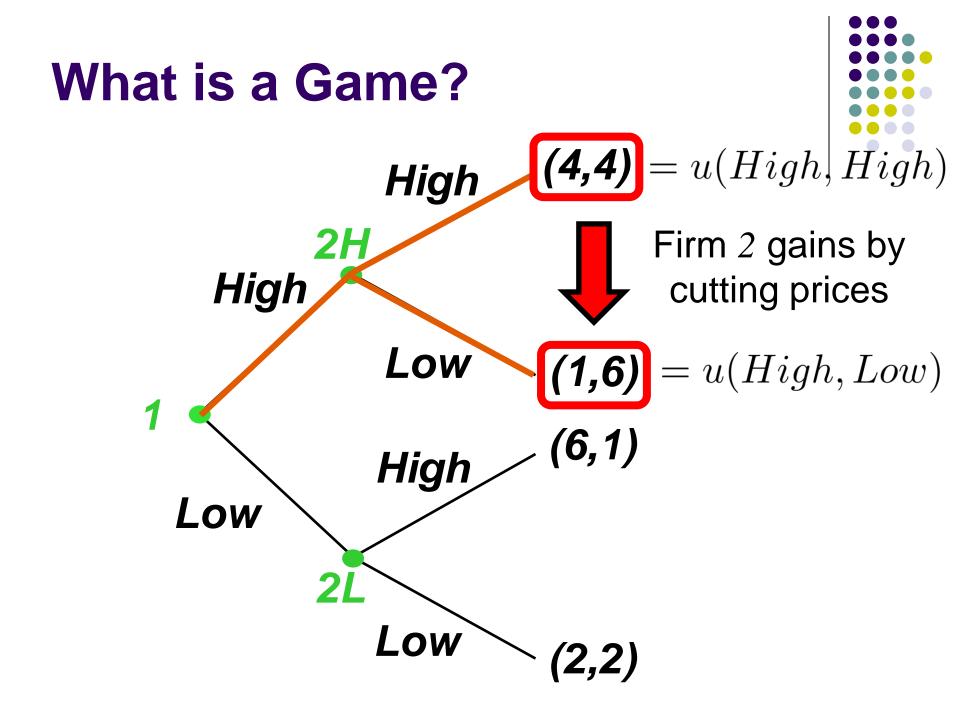


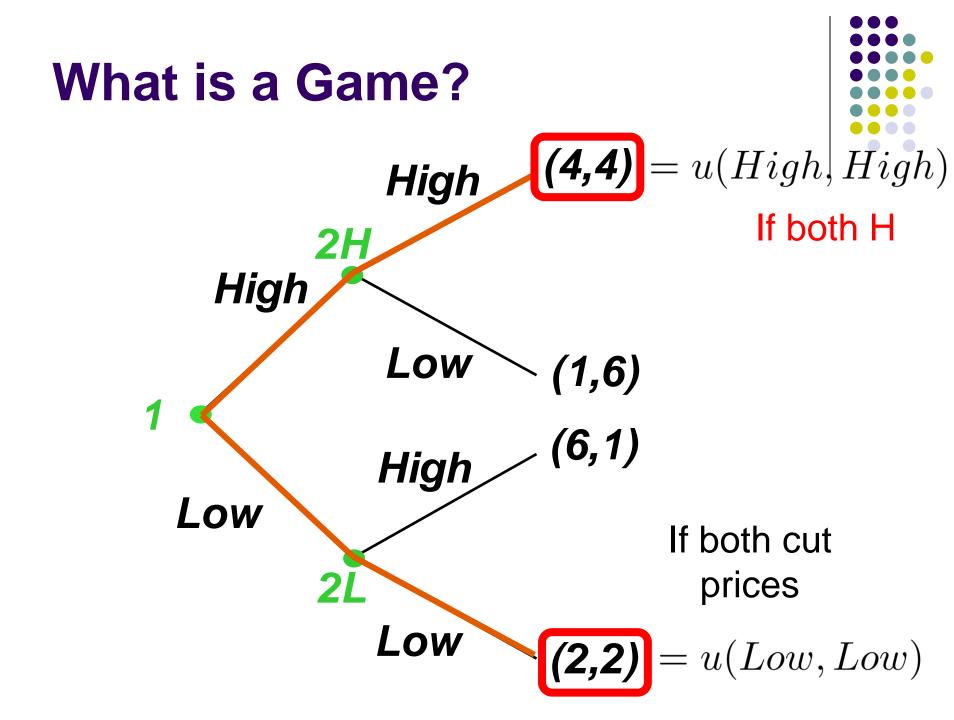
- Example: Two competing firms
- Agents i = manager of firm i = 1, 2
- Post next week's price on Sunday Times
 - High price or Low price
- Agent 1 sets price first
 - Sunday Times posts price online instantly; Agent 2 sees opponent's price before setting own price
- Represent game as a game tree

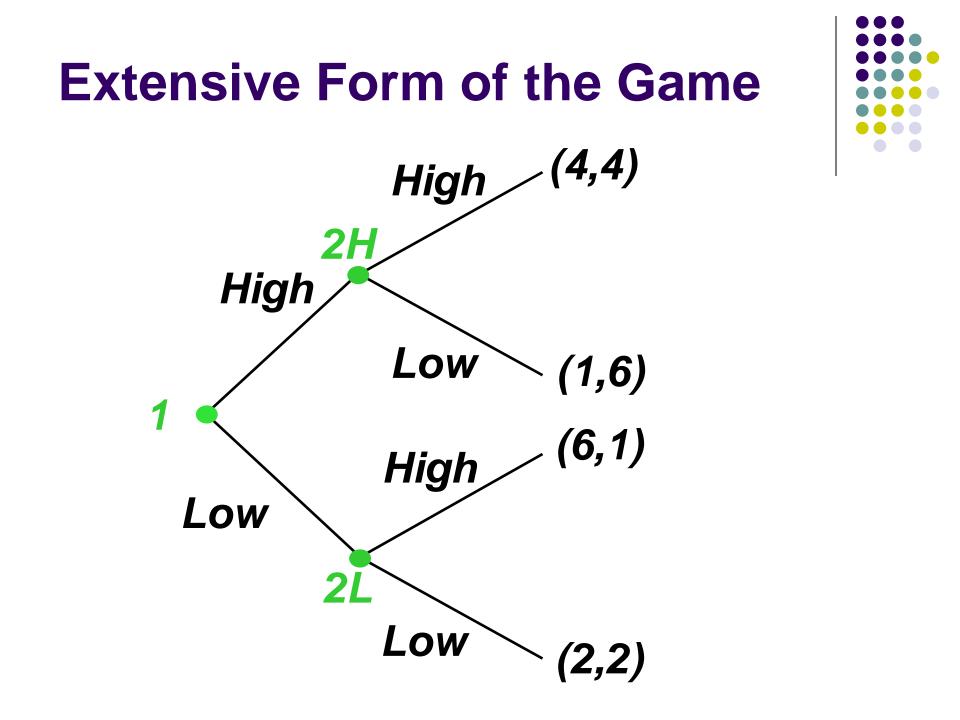




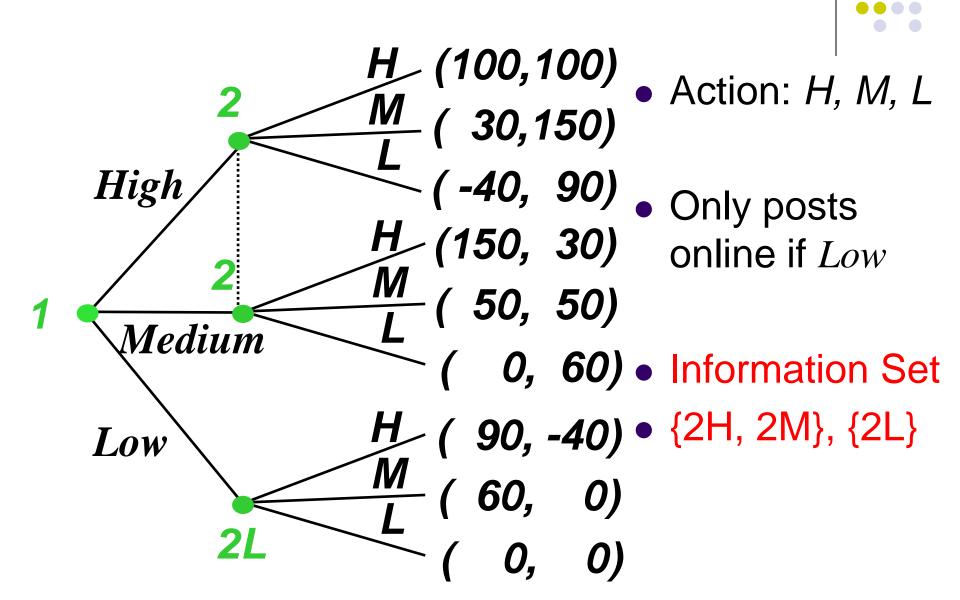


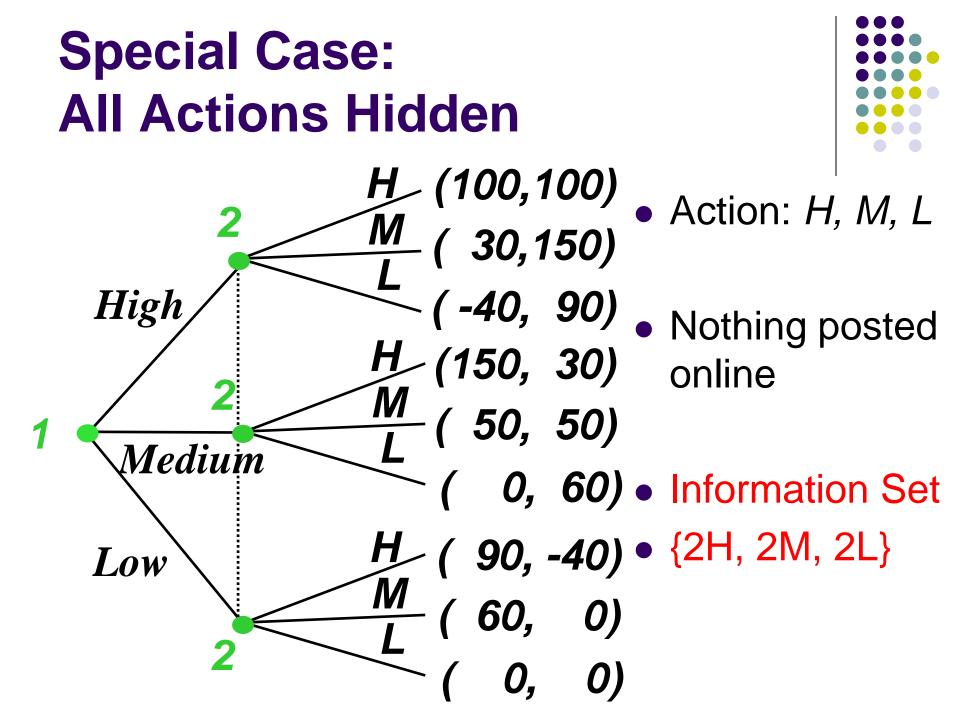






Other Extensive Form Games



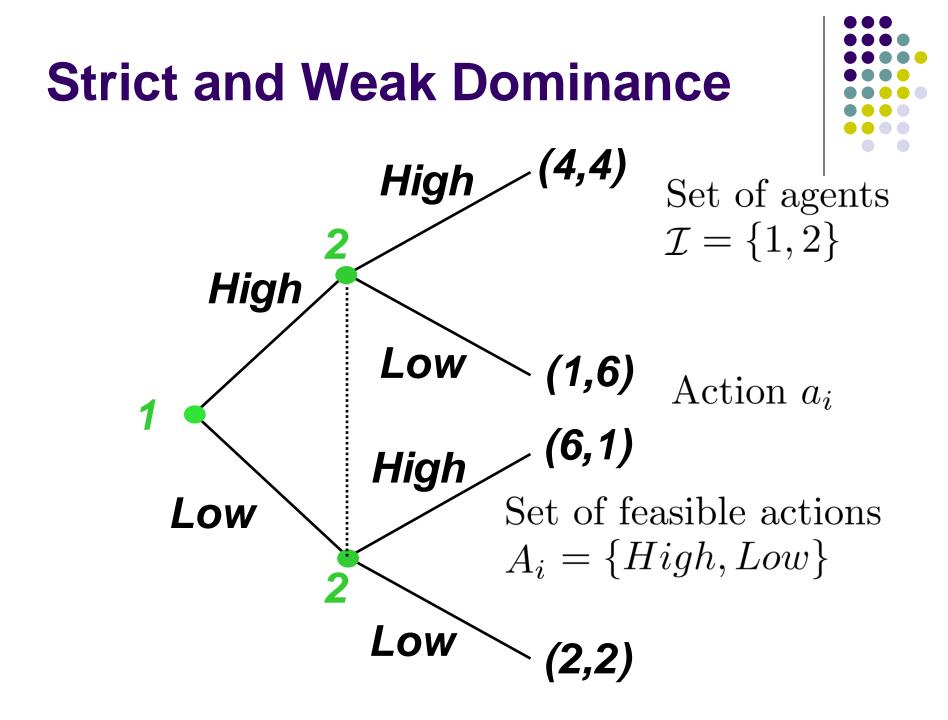


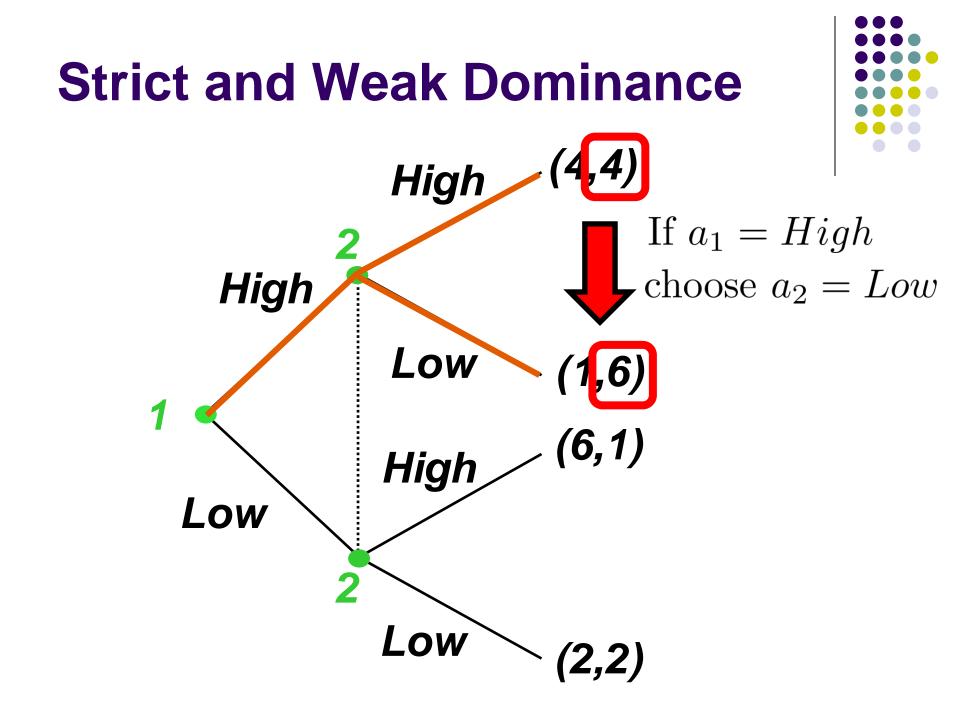
Strict and Weak Dominace

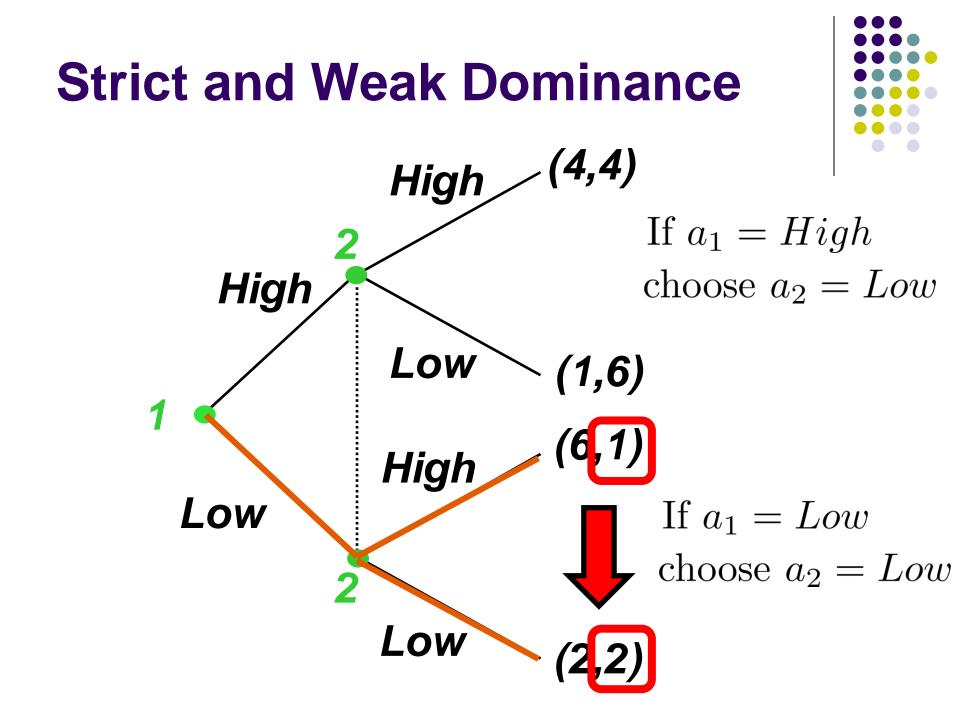
- Set of opponent action space $A_{-i} = \bigotimes$ A_i $j \neq i$
- For agent *i*,
 - a_i is strictly dominated by $\overline{a_i}$ if

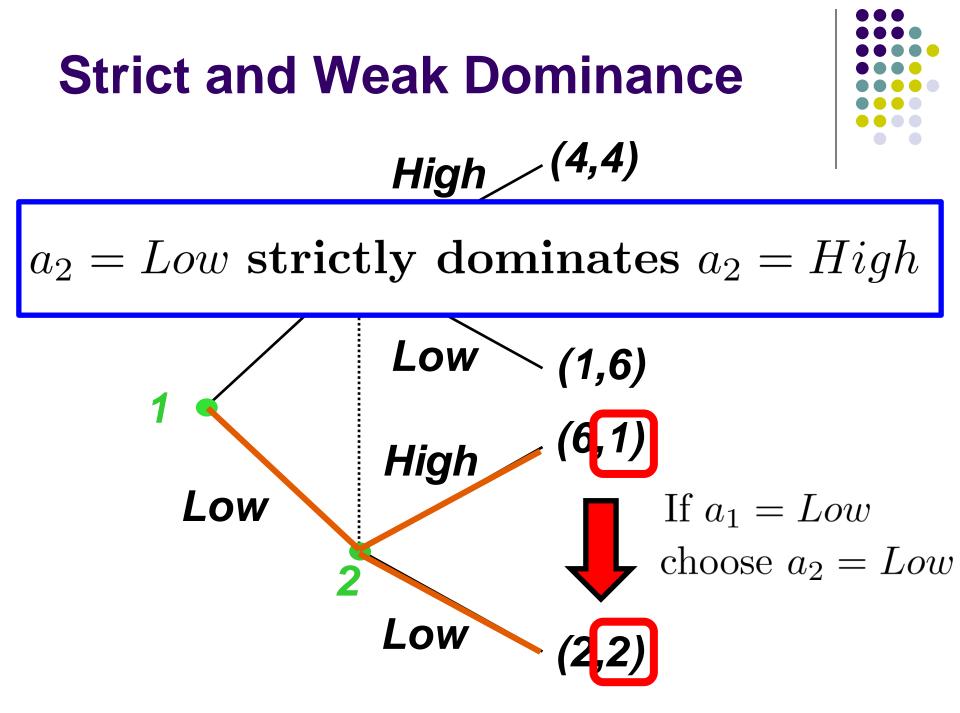
 $u_i(\overline{a_i}, a_{-i}) > u_i(a_i, a_{-i})$ for all $a_{-i} \in A_{-i}$

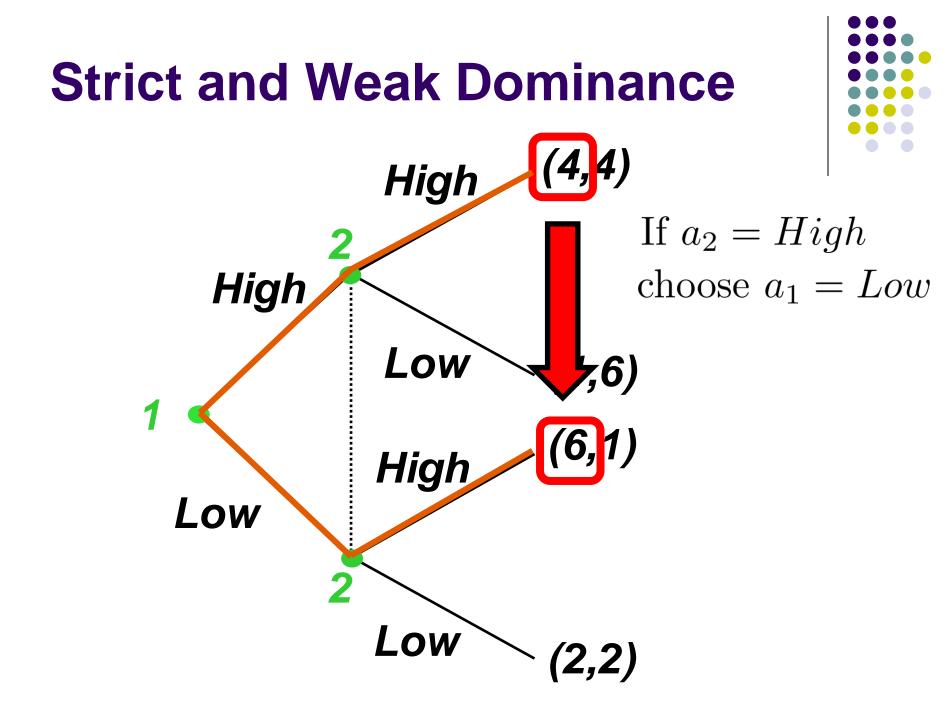
 a_i is weakly dominated by $\overline{a_i}$ if $u_i(\overline{a_i}, a_{-i}) \ge u_i(a_i, a_{-i})$ for all $a_{-i} \in A_{-i}$ $u_i(\overline{a_i}, a_{-i}) > u_i(a_i, a_{-i})$ for some $a_{-i} \in A_{-i}$

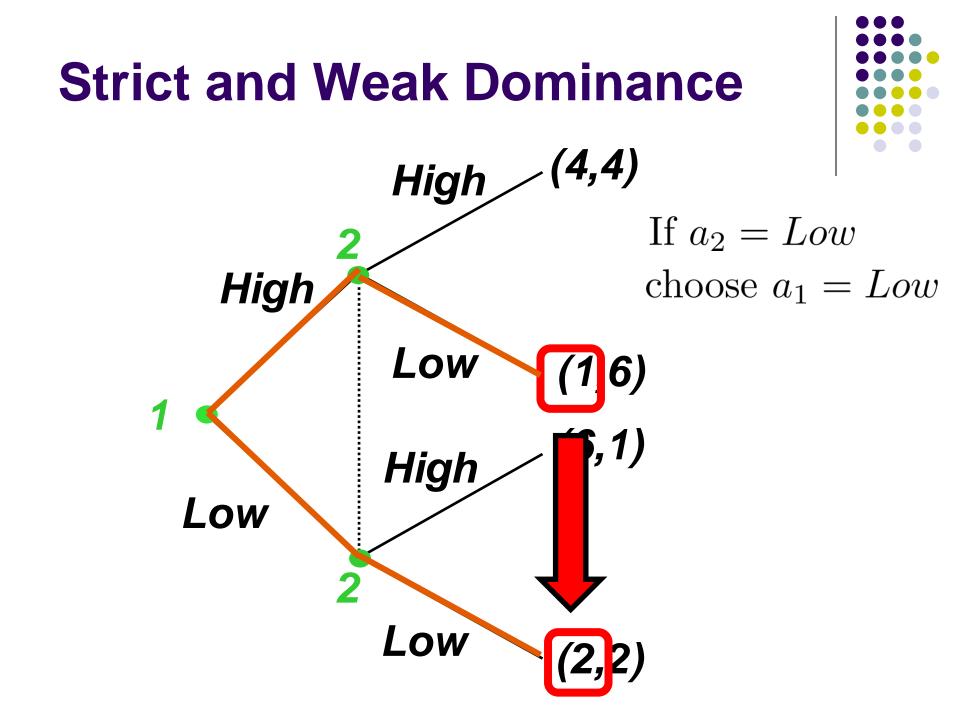


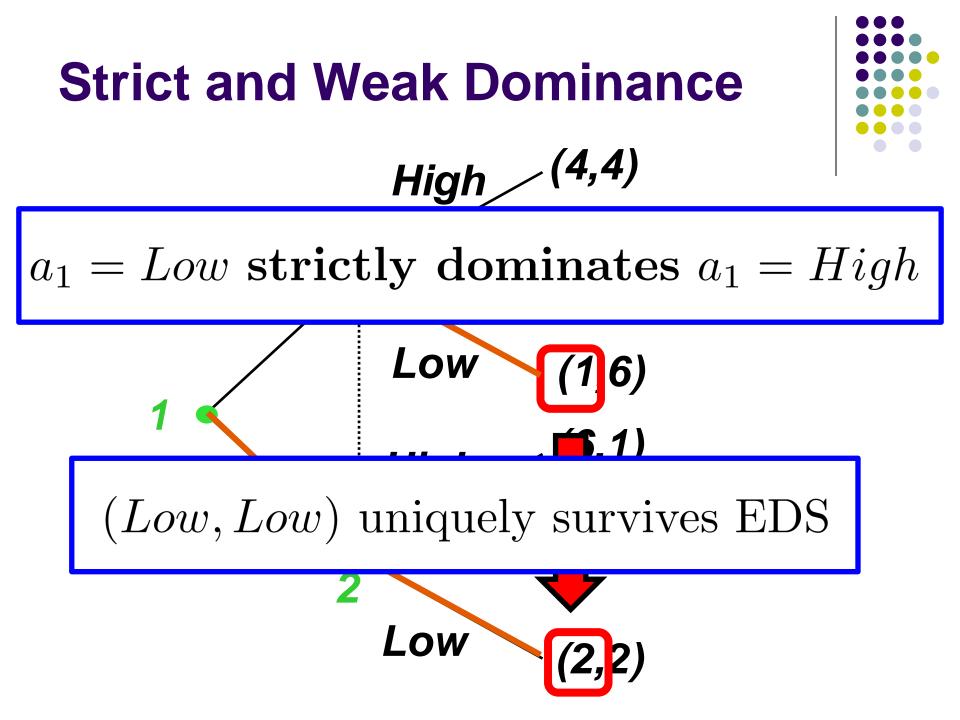






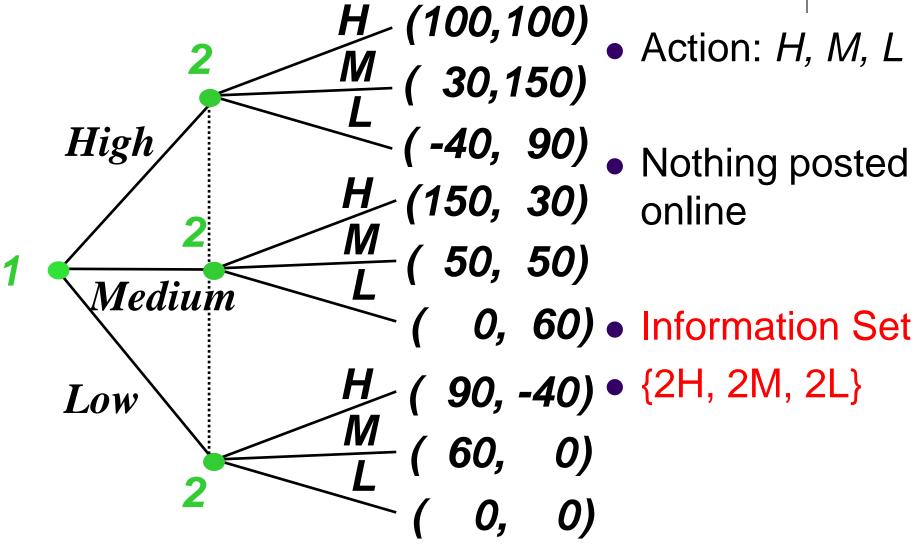






Simultaneous Game: Extensive Form





Simultaneous Game: Strategic Form (Normal Form)



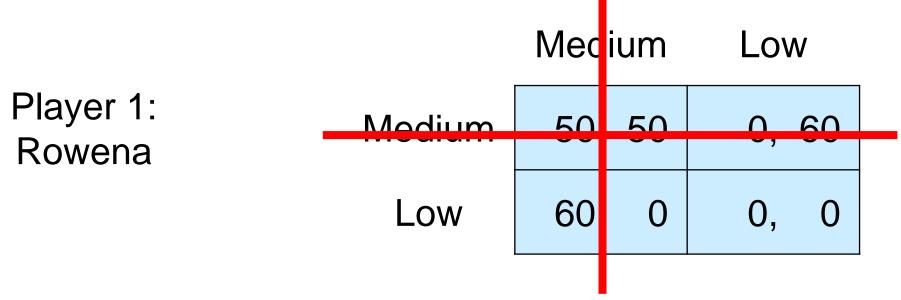


Elimination of Dominated Strategies (EDS)



Medium weakly dominated by *Low*

Player 2: Colin



Iterative Elimination of Dominated Strategies



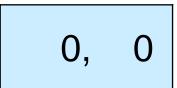
Player 2: Colin

(Low, Low) uniquely survives IEDS

Low

Player 1: Rowena

Low

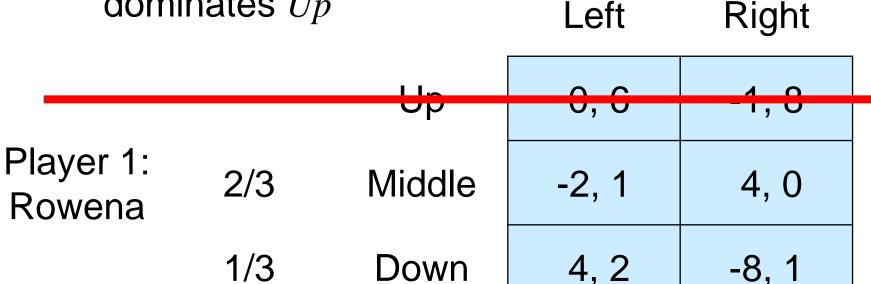


Mixed Strategy and Dominance



(2/3,1/3)-mixture of (Middle, Down) weakly dominates Up

Player 2: Colin



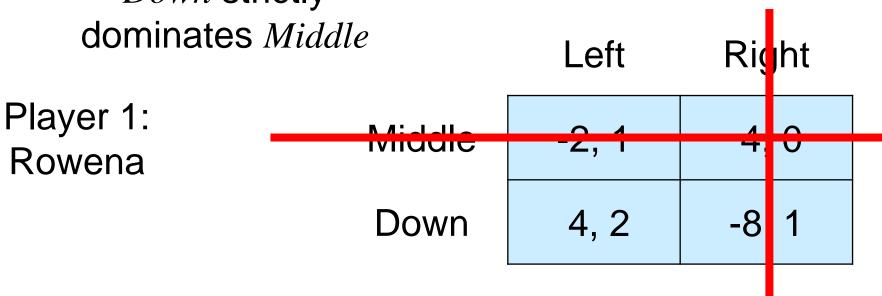
Mixed Strategy and IEDS



Left strictly dominates *Right*

Down strictly





Equilibrium of "One-Shot" Simultaneous Game

- Each Agent $i \in \mathcal{I}$
- Has finite Action Set $A_i = \{a_{i1}, a_{i2}, \cdots, a_{im}\}$
- Agent *i*'s Strategy Set $S_i = \Delta(A_i) = \left\{ \pi \middle| \pi \ge 0, \sum_{j=1}^{m_i} \pi_j = 1 \right\}$
- Mixed Strategy: $\pi_i(a_i)$
- Strategy Profile:

$$s = (s_1, \cdots, s_I) \in S = S_1 \times \cdots \times S_I$$

Equilibrium of "One-Shot" Simultaneous Game

- Consequence of the game (for agent *i*): $\pi_i(a)$
- Outcome of the game (for agent *i*): $x_i(a)$
- Agent *i*'s Expected Utility

$$u_i = \sum_{a \in A} \pi_i(a) v_i(x_i(a)) = u_i(a) \cdot \pi_i(a)$$

- Mixing in Continuous Action Space: $\mu_i \in \Delta(A_i)$
- Expected Utility in Continuous Action Space:

$$u_i(s) = \int_{a \in A} u_i(a) d\mu(a)$$



Nash Equilibrium



- Strategy Profile: $s \in S = \Delta_1(A_1) \times \cdots \times \Delta_I(A_I)$
- Best Response: $BR_i(s_{-i})$
- Best Response Mapping: $BR(s) = (BR_1(s_{-1}), \cdots, BR_I(s_{-I}))$
- Nash Equilibrium: s such that BR(s) = s
 - Fixed Point in the BR mapping
- Consider a strategy profile $\overline{s} = (\overline{s}_1, \cdots, \overline{s}_I)$
- Is there any other strategy strictly better for agent *i* (if others play according to \overline{s}_{-i})

Nash Equilibrium



- For simultaneous game played by agents 1~I
- The strategy profile $\overline{s} = (\overline{s}_1, \dots, \overline{s}_I)$ is a Nash Equilibrium if the strategies are mutual BR.
- In other words,
- For each $i \in \mathcal{I}$ and all $a_i \in A_i$ $u_i(\overline{s}_i, \overline{s}_{-i}) \ge u_i(a_i, \overline{s}_{-i})$
 - Note that you only need to check pure strategies since mixed strategies yield a weighted average of payoffs among pure strategies

Nash Equilibrium: Partnership Game



- Two Agents have equal share in a partnership
- Choose Effort: $a_i \in A_i = \{1, 2, 3\}$
- Total revenue: $R = 12a_1a_2$
- Cost to agent *i*: $C_i(a_i) = a_i^3$
- Payoff: $u_i(s) = R C_i(a_i) = 12a_1a_2 a_i^3$
- Game matrix and Nash Equilibrium...

30

Nash Equilibrium: Partnership Game

1 is a BR if other picks 1 Player 2: Colin 2 is a BR if other 2 3 picks 2 or 3 11, 4 17, -9 5, 5 1 Player 1: 2 4, 11 <u>16</u>, 16 <u>28</u>, 9 Rowena 3 9, 28 27, 27 -9, 17

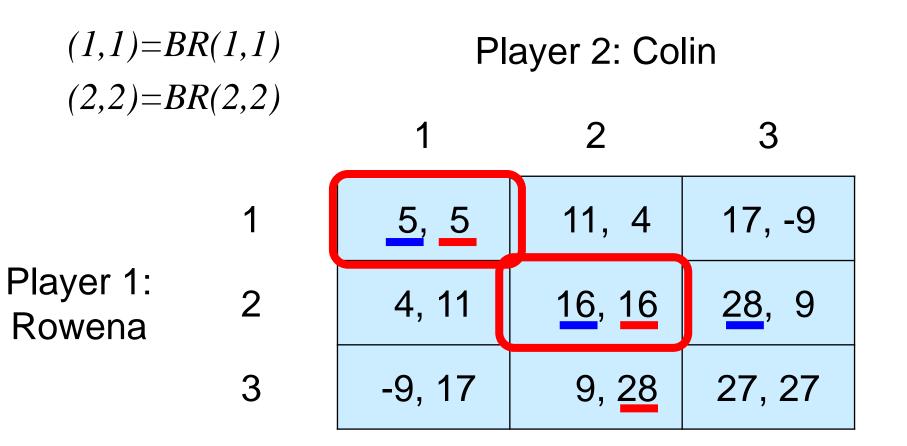


Nash Equilibrium: Partnership Game

- 1 is a BR if other picks 1 2 is a BR if other picks 2 or 3 1 Player 1: 2 Rowena 3
 - Player 2: Colin 2 3 11, 4 17, -9 5, 5 <u>16, 16</u> 28, 4, 11 9 -9, 17 9, <u>28</u> 27, 27







Nash Equilibrium: Partnership Game

- This is NOT the only two NE
- Solve for MSE:

• For
$$s_2 = (p, 1 - p, 0) \in \Delta(A_2)$$

 $u_1(1, s_2) = 5p + 11(1 - p) = 11 - 6p$
• $= u_1(2, s_2) = 4p + 16(1 - p) = 16 - 12p$

• Hence, $p = \frac{5}{6}$ • By symmetry, MSE is $s_1 = s_2 = \left(\frac{5}{6}, \frac{1}{6}, 0\right)$



Common Knowledge

- Common Knowledge of the Game
- Common Knowledge of Rationality
- Common Knowledge of Equilibrium
- Exercise: Is "九二共識" truly a consensus in terms of common knowledge?



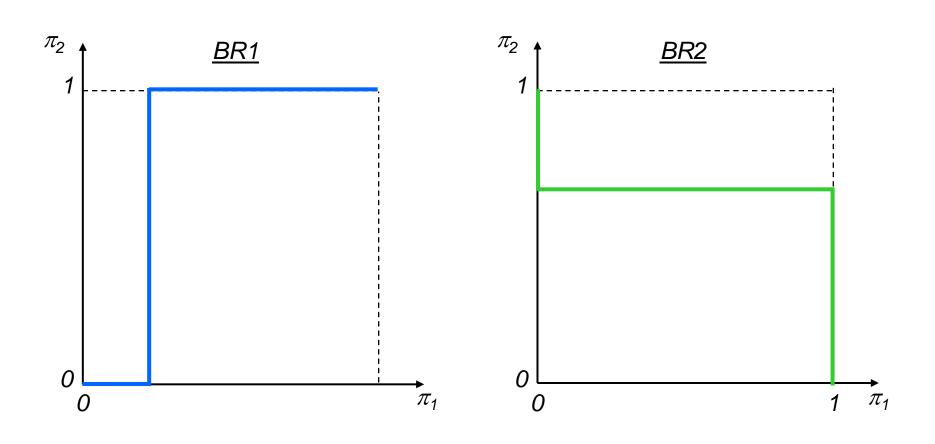


- Use: Kakutani's Fixed Point Theorem (FPT) If S ⊆ Rⁿ is closed, bounded & convex and if φ is an upper hemi-continuous correspondence from S to S, such that φ(s) is non-empty and convex, then φ(s) has a fixed point.
- Proposition 9.1-1: Existence of NE (Nash, 1950)
- In a game with finite action sets, if players can choose either pure or mixed strategies, there exists a Nash Equilibrium.

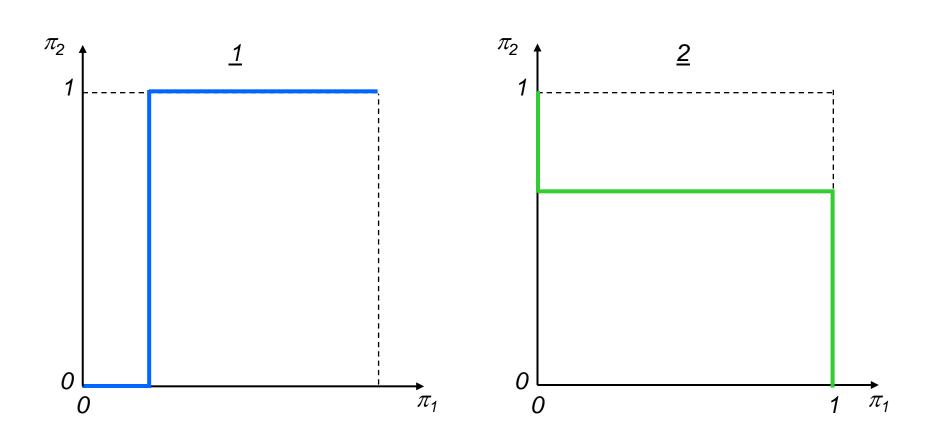


- Consider the following simpler FPT:
 If S₁, S₂ ⊆ R is closed, bounded and convex and φ₁(s₂), φ₂(s₁) are continuous functions from S_{-i} to S_i, then φ = (φ₁, φ₂) has a fixed point.
- Existence of Nash Equilibrium requires:
- Strategy sets are closed, bounded and convex,
- BR functions are indeed continuous...



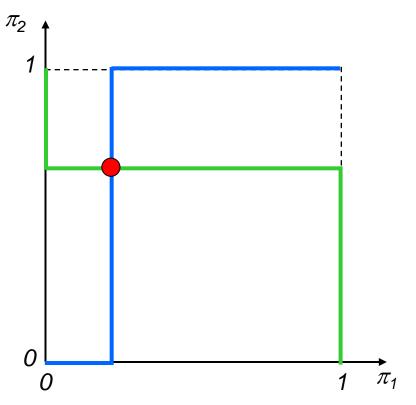








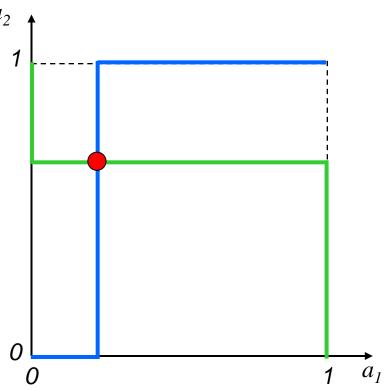
Mixed-strategy NE in which player <u>1</u> plays Up with probability π_1 and player <u>2</u> plays Left with probability π_2 .



Existence of Equilibrium: For Continuous Action Space



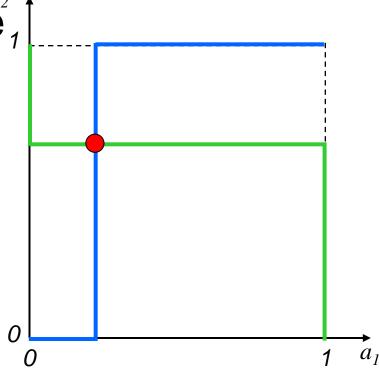
For continuous action space a_2 (where each player chooses a pure strategy a_i), there exists a pure strategy NE in which player <u>1</u> plays a_1 and player 2 plays a_2 .



Existence of Equilibrium: For Non-unique BR

- Why do we need Kakutani's FPT?
- Because best response₁
 may not be unique!!!
- BR correspondences,
 - Not only BR "functions"
- Upper hemi-continuous
 - Not "Continuous"







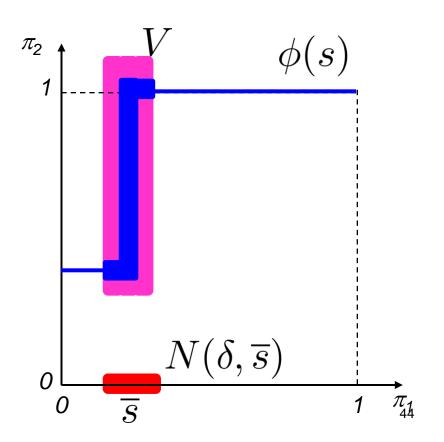
- Use: Kakutani's Fixed Point Theorem (FPT) If S ⊆ Rⁿ is closed, bounded & convex and if φ is an upper hemi-continuous correspondence from S to S, such that φ(s) is non-empty and convex, then φ(s) has a fixed point.
- Closed and Bounded
- Convex
- Upper hemi-continuous



- \bullet Bounded $S\subseteq B(s,r), r<\infty$
 - Contained in a ball of radius r (centered at s)

• Convex If
$$s^0, s^1 \in C$$
, for $0 < \lambda < 1$,
 $s^{\lambda} = (1 - \lambda)s^0 + \lambda s^1 \in C$.

- $\phi(s)$ is upper hemicontinuous at \overline{s} if
- For any open neighborhood $V ext{ of } \phi(\overline{s})$
- There exists $N(\delta, \overline{s})$ a δ -neighborhood of \overline{s}
- such that $\phi(s) \subseteq V$ for all $s \in N(\delta, \overline{s})$







- Using Kakutani's Fixed Point Theorem (FPT)
- Proposition 9.1-1: Existence of NE (Nash, 1950)
- In a game with finite action sets, if players can choose either pure or mixed strategies,
 - Mixed strategy profile ($\pi_1, \pi_2, ..., \pi_n$), $0 \leq \pi_i \leq 1$
 - Closed, bounded and convex
- there exists a Nash Equilibrium.
 - BR correspondence is non-empty, convex (mixing among BR is also BR), and upper hemi-continuous



- Proposition 9.1-2: Existence of pure NE
- In a game with action sets $A_i \subset \mathbf{R}^n$ is closed, bounded and convex, and utility *u* is continuous,
- If BR sets $BR_i(a_{-i}) \subseteq A_i$ are convex,
- there exists a pure strategy Nash Equilibrium.
- Corollary 9.1-3: Existence of pure NE
- If BR sets $BR_i(a_{-i}) \subseteq A_i$ are single-valued, or If $u_i(a_i, a_{-i})$ are quasi-concave over a_i
- there exists a pure strategy Nash Equilibrium.

Summary of 9.1

- Game Tree
 - Extensive Form and Information Sets
- Simultaneous Game
 - Strategic Form (Normal Form)
- Nash Equilibrium
 - Existence of Nash Equilibrium (by Kakutani's FPT)
- HW 9.1: Riley 9.1-1~4